

MASSACHUSETTS INSTITUTE OF TECHNOLOGY
SLOAN SCHOOL OF MANAGEMENT

15.565 Integrating eSystems:

Technology, Strategy, and Organizational Factors

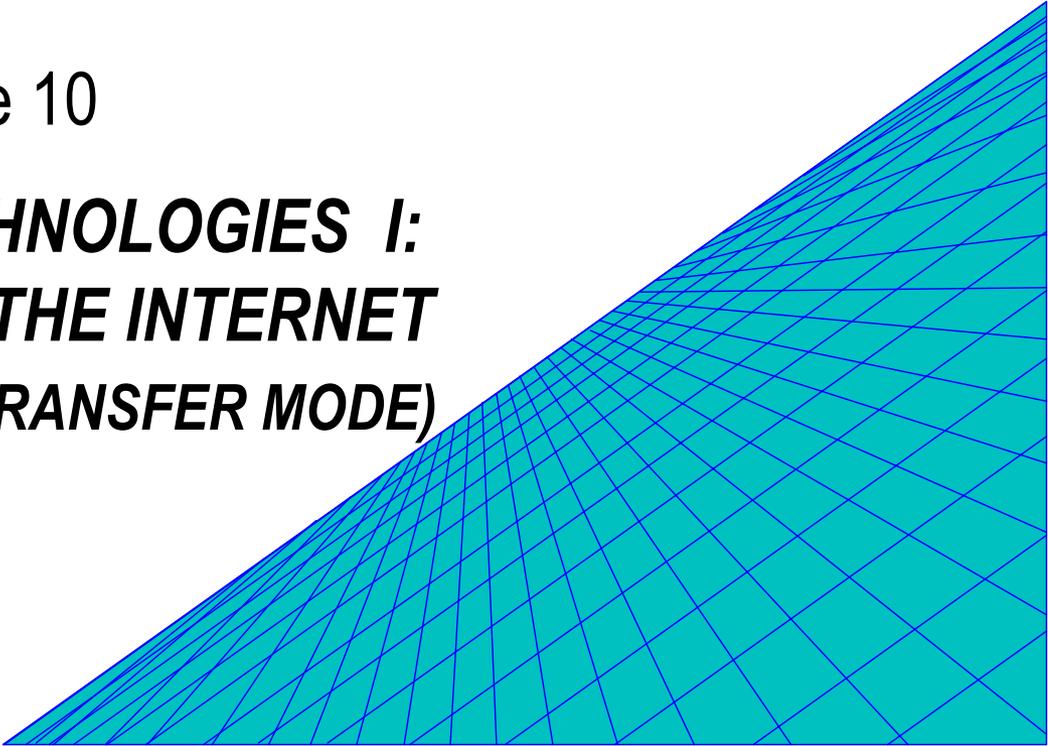
15.578 Global Information Systems:

Communications & Connectivity Among Information Systems

Spring 2002

Lecture 10

***EMERGING TECHNOLOGIES I:
THE FUTURE OF THE INTERNET
(ASYNCHRONOUS TRANSFER MODE)***



PACKET ROUTING ISSUES

• PHILOSOPHIES

- **DATAGRAM**: EACH PACKET INDEPENDENT
 - LOAD SPREAD
- **VIRTUAL CIRCUIT**: ALL PACKETS OF A MESSAGE USE SAME PATH
 - RESERVE RESOURCES & ARRIVE IN ORDER
 - MORE “SET UP” OVERHEAD

• ROUTING

- EFFICIENT PATH & CONGESTION CONTROL
- MANY APPROACHES

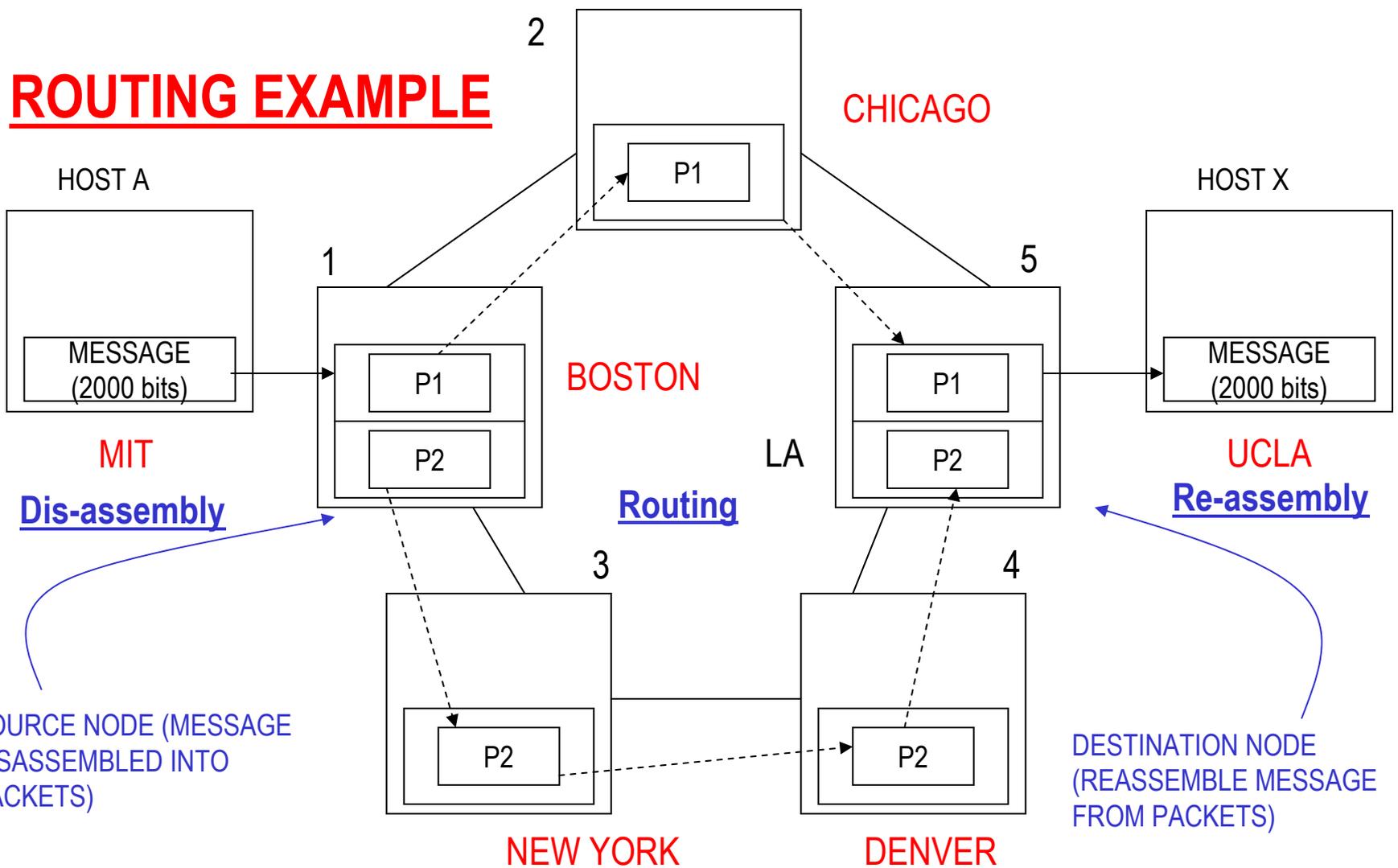
• ROUTING ALGORITHMS

- STATIC DIRECTORY
- DYNAMIC ROUTING
- DECENTRALIZED SCHEMES

• CONGESTION CONTROL

- WANT TO AVOID “GRIDLOCK”
- APPROACHES: PERMITS, “SLOW DOWN” REQUEST, DISCARD

ROUTING EXAMPLE



ADVANTAGES:

- NOT TIE UP COMMUNICATION LINES FOR LONG PERIODS (I.E., NO LONG MESSAGES) -- INTERLEAVE PACKETS
- ON COMMUNICATION ERROR, ONLY NEED TO RETRANSMIT PACKET AFFECTED RATHER THAN ENTIRE MESSAGE
- BUFFER MANAGEMENT SIMPLER / LESS DEADLY EMBRACE
- ABLE TO SHIP PACKETS IN PARALLEL (**PARALLEL PROCESS**)
- ABLE TO SEND PARTIAL MESSAGE (PACKET) WITHOUT WAITING TO RECEIVE ENTIRE MESSAGE (**PIPELINE PROCESS**)

THE COLLAPSE (IMPLOSION) OF “POTS”

- AT&T LONG-DISTANCE NETWORK
 - HANDLES 125M CALLS/DAY
 - AVERAGE CALL = 5 MINUTES
 - EACH CALL USES 64K bps
 - TOTAL = 28.8G bps (AVERAGED)
 - MUCH LESS THAN CAPACITY OF SINGLE OPTICAL FIBER !
- EVEN IF TOTAL POPULATION TALKED 24 HR/DAY
 - TOTAL = 12.8T bps
- TREMENDOUS EXCESS CAPACITY
 - “DARK FIBER”
- WHY DOESN'T / HASN'T “POTS” COLLAPSED?

QUALITY OF SERVICE (QoS) TRADEOFFS / REQUIREMENTS

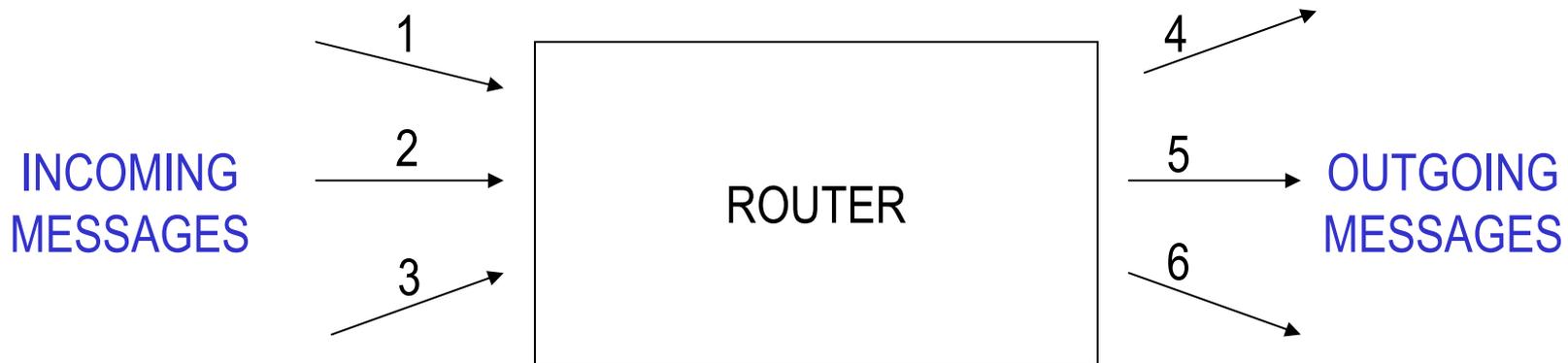
- INTERNET GIVES “BEST EFFORT”, BUT NO GUARANTEE

RELIABILITY

		LOW	HIGH
<u>TIMELINESS</u>	LOW		DATA (INTERNET) [PACKET SWITCH]
	HIGH	VOICE ("POTS") [CIRCUIT SWITCH]	

- CAN A SINGLE NETWORK SERVICE BOTH NEEDS?
- OTHER DIFFERENCES (TERMINOLOGY, REGULATIONS, DEPRECIATION, RATE-OF-CHANGE)
- PLUS OTHER APPLICATION NEEDS (E.G., BULK FILES, MOVIES, HDTV, ETC.)

CAUSE OF DELAYS IN PACKET NETWORK



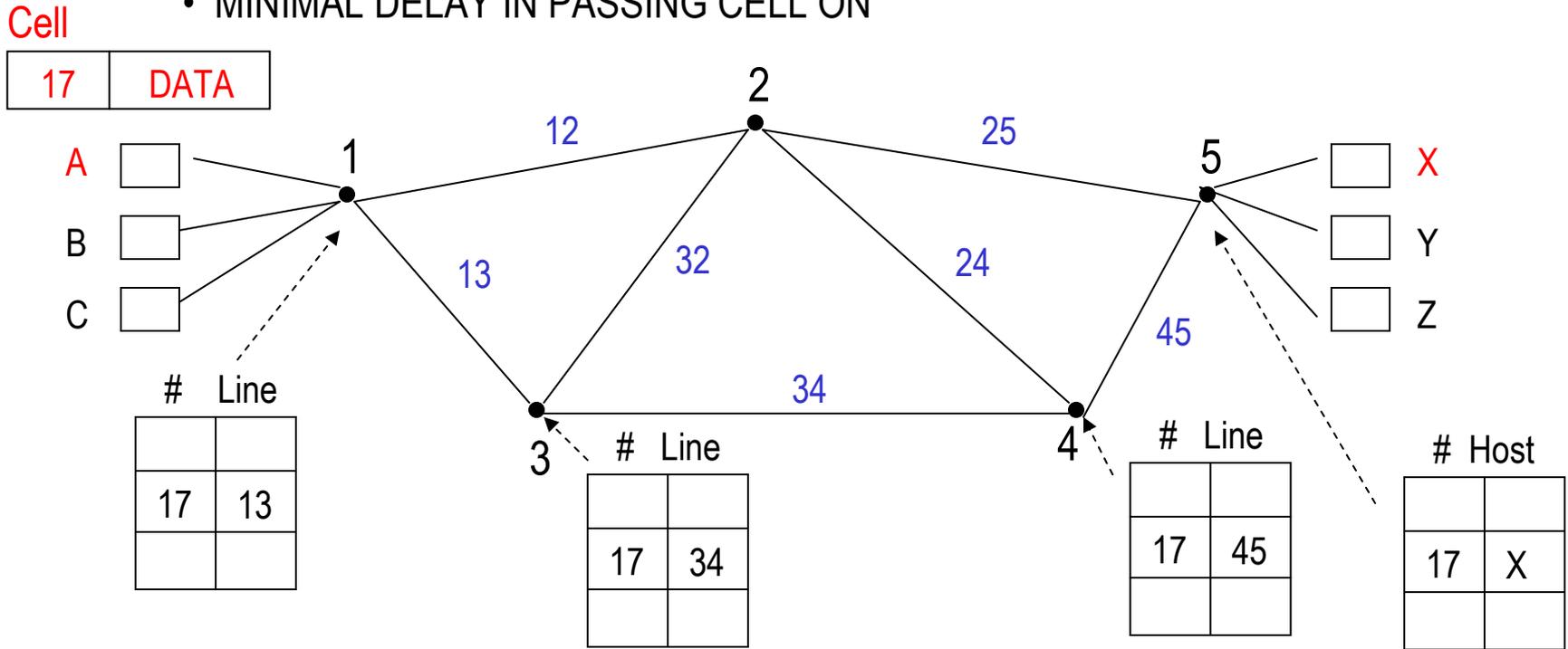
- **PACKET SIZE AND BACKLOG**
 - CAN'T SEND NEXT PACKET, UNTIL CURRENT ONE SENT
 - MAY ALREADY BE SEVERAL WAITING TO BE SENT OVER SAME LINE
- **ROUTER SPEED**
 - HOW FAST DETERMINE OUTGOING LINE
 - OTHER FACTORS: PRIORITY, ERROR DETECTION
- **ALTERNATE ROUTE IMPACTS**
 - PACKETS MAY ARRIVE OUT OF ORDER
 - MUST WAIT FOR CORRECT SEQUENCE TO BE RE-ASSEMBLED

PACKET SIZE SOLUTION (AND PROBLEM !)

- ASYNCHRONOUS TRANSFER MODE (ATM)
 - ASYNCHRONOUS (E.G., PACKETS) VS. SYNCHRONOUS (E.G., TDM)
- STANDARDIZED 53 BYTE PACKET - - WHY 53?
 - 5 BYTE HEADER
 - 48 BYTES DATA
- IMPACT ON ROUTER (BASICALLY A SPECIALIZED COMPUTER)
 - PREVIOUS:
 - 1000 BYTE PACKET @ 1 M bps = 10 ms (1/100th of sec)
 - FOR A 100 MIP NODE: 10 ms = 1,000,000 INSTRUCTIONS
 - IF ROUTING TOOK 10,000 INSTRUCTIONS = 1% ROUTER CAPACITY
 - NOW:
 - 53 BYTE PACKET @ 100 M bps = 5 μ S (5 millionths of a sec)
 - FOR A 1000 MIP NODE, 5 μ S = 5,000 INSTRUCTIONS
 - IF ROUTING STILL TOOK 10,000 INSTRUCTIONS = 200% OF CAPACITY !
 - FUTURE:
 - HOW HANDLE 1G+ bps LINES?

ATM PATH ROUTING STRATEGY (Virtual Circuit)

- KEY CONCEPT: PRE-CALCULATE AND PRE-ALLOCATE PATH
- NO SOURCE OR DESTINATION ADDRESSES USED IN CELL
- MINIMAL DELAY IN PASSING CELL ON



#	Line	New #
17	13	63

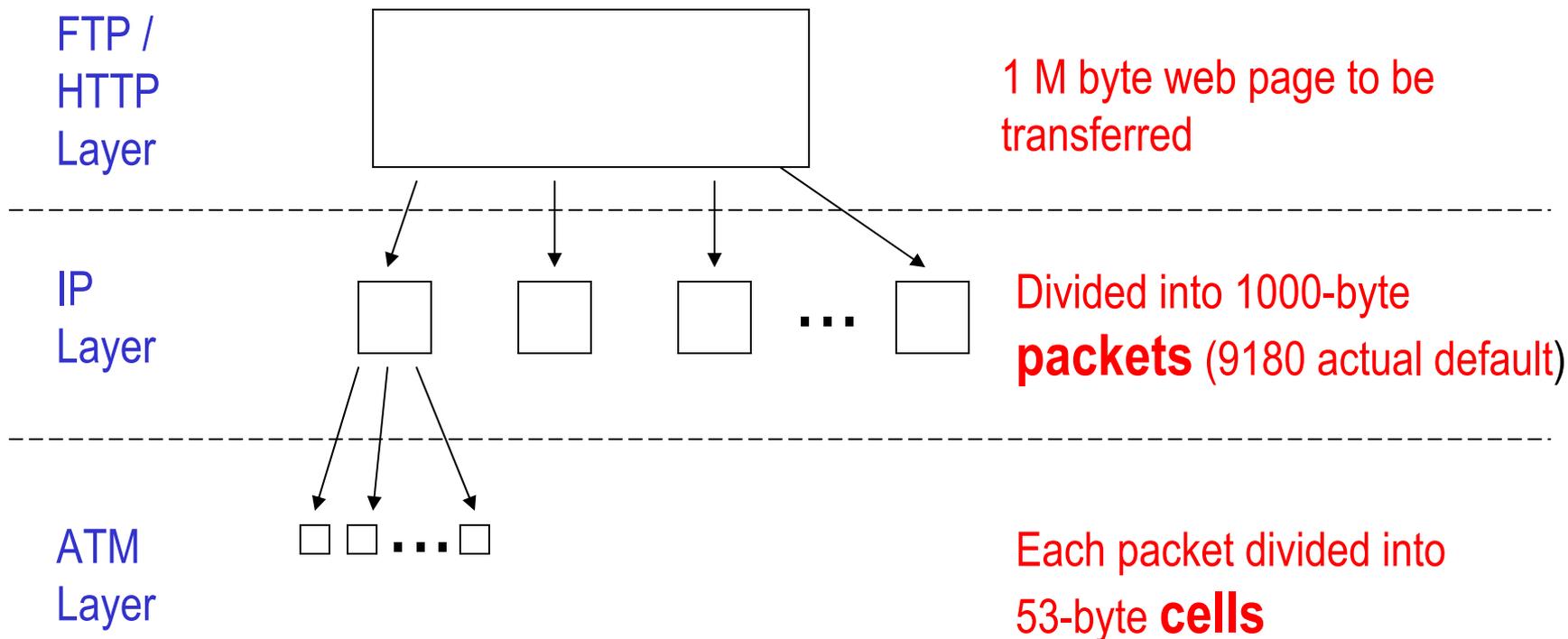
#	Line	New #
63	34	75

#	Line	New #
75	45	12

#	Host
12	X

RELATIONSHIP OF IP TO ATM

- ATM ADAPTATION LAYERS (AAL5) [See Appendix for more details]



- ADDRESS MAPPING:

- MUST MAP IP ADDRESS TO ATM ADDRESS AT “SETUP”

- USUALLY VIA A SHARED SERVER (WITH KNOWN ATM ADDRESS)

- MUST MAP EACH IP ADDRESS TO ATM FOR EACH PACKET AFTER SETUP

ATM CHALLENGES / OBSTACLES

- ATM HARDWARE (CURRENTLY) EXPENSIVE
- MIGRATION COSTS AND COMPLEXITIES
- STANDARDS BATTLES/COMPROMISES
 - MULTIPLE SELF-INTERESTS
 - E.G., ATM ADDRESS FORMAT DISPUTE (BOTH SUPPORTED)
 - TELEPHONE: 8-OCTET FORMAT (ISDN PHONE NUMBER)
 - CCITT
 - ATM FORUM: 20-OCTET NETWORK SERVICES ACCESS POINT (NSAP)
 - CISCO, IBM, etc.

FUTURE OF TCP/IP (IPng, IPv6)

- CURRENT IPv4 (SINCE 1970s) LIMITATIONS
 - IP ADDRESS = 32 BITS (MAX 4 BILLION NODES)
 - REAL-TIME AUDIO AND VIDEO POOR
- MAINTAIN
 - CONNECTIONLESS DELIVERY (EACH DATAGRAM INDEPENDENT)
 - SENDER CHOOSES DATAGRAM SIZE (MAX = 64K BYTES)
 - SENDER SPECIFIES MAXIMUM HOPS
- NEW (IPv6)
 - ADDRESS = 128 BITS ($2^{128} = 3.4 \times 10^{38}$) - WHY?
 - PRE-ALLOCATION OF NETWORK RESOURCES (“FLOW”)
 - RESERVE AND RESTRICT
 - MAXIMUM FRAGMENT SIZE DETERMINED IN ADVANCE
 - NOT AS DYNAMIC
 - PROVISION FOR PROTOCOL EXTENSION
 - MORE FLEXIBLE HEADER FORMAT/OPTIONS
 - SUPPORT: UNICAST, CLUSTER, MULTICAST

IPv6 ADDRESS ISSUES

- FORMATS:

- TRADITIONAL IPv4 = 104. 230. 140. 100
- IPv6 = 104. 230. 140. 100. 255. 255. 255. 255. 0. 0. 17. 128. 150. 10. 255. 255
- IPv6 = 68E6: 8C64: FFFF: FFFF: 0: 1180: 96A: FFFF (HEXADECIMAL)
- ALSO PROVISIONS FOR ABBREVIATIONS, ESP IF SEQUENCES OF 0's (::)

- ADDRESS SPACE ALLOCATION

- TRADITIONAL IPv4 - 2 PARTS:
 - NETWORK PREFIX (CENTRAL MANAGED)
 - HOST SUFFIX (ASSIGNED BY ORGANIZATION)
- IPv6 OFFER VARIES POSSIBILITIES, SUCH AS:
 - PROVIDER ID
 - SUBSCRIBER ID
 - SUBNET ID
 - NODE ID (E.G., IF 48 BITS CAN = ETHERNET ADDRESS)

OTHER INTERNET EMERGING TECHNOLOGIES & ISSUES

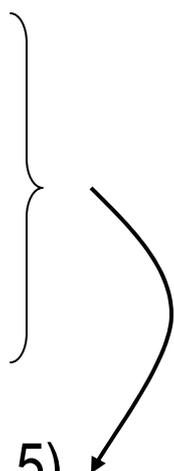
- TECHNOLOGIES
 - SATELLITE (POINT-TO-POINT, SYNCHRONIZED)
 - PACKET RADIO
 - etc.
 - MULTI-MEDIA
 - REAL-TIME AUDIO AND VIDEO
-

EXPERIMENTS

- NEXT GENERATION INTERNET (NGI)
 - DARPA FUNDED
 - 1G bps → 1 T bps (OVER ATM)
- INTERNET II (www.internet2.org -- *not .com*)
 - 170 UNIVERSITIES
 - 622 M bps → 2.4 G bps
- OTHER EXPERIMENTS

APPENDIX: ATM ADAPTION LAYERS

- TYPE 1 - CONSTANT BIT RATE (CBR)
 - VOICE AND VIDEO
 - TYPE 2 - VARIABLE BIT RATE (VBR)
 - PACKET VIDEO
 - TYPE 3 - CONNECTION-ORIENTED VBR DATA TRANSFER
 - FILE TRANSFERS
 - TYPE 4 - CONNECTIONLESS VBR DATA TRANSFER
 - SHORT BURSTY
 - TYPE 5 - SIMPLE & EFFICIENT ADAPTION LAYER (SEAL/AAL5)
 - TYPE 3 WITHOUT ERROR DETECTION

(REPLACES 3 AND 4)
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ATM NETWORK INTERFACES

