

15.083J/6.859J Integer Optimization

Lecture 9: Duality II

## 1 Outline

SLIDE 1

- Solution of Lagrangean dual
- Geometry and strength of the Lagrangean dual

## 2 The TSP

SLIDE 2

$$\begin{aligned}
 & \sum_{e \in \delta(\{i\})} x_e = 2, \quad i \in V, \\
 & \sum_{e \in E(S)} x_e \leq |S| - 1, \quad S \subset V, S \neq \emptyset, V, \\
 & x_e \in \{0, 1\}. \\
 \text{min} \quad & \sum_{e \in E} c_e x_e \\
 \text{s.t.} \quad & \sum_{e \in \delta(\{i\})} x_e = 2, \quad i \in V \setminus \{1\}, \\
 & \sum_{e \in \delta(\{1\})} x_e = 2, \\
 & \sum_{e \in E(S)} x_e \leq |S| - 1, \quad S \subset V \setminus \{1\}, S \neq \emptyset, V \setminus \{1\}, \\
 & \sum_{e \in E(V \setminus \{1\})} x_e = |V| - 2, \\
 & x_e \in \{0, 1\}.
 \end{aligned}$$

Dualize  $\sum_{e \in \delta(\{i\})} x_e = 2, i \in V \setminus \{1\}$ .  
What is the relation of  $Z_D$  and  $Z_{LP}$ ?

## 3 Solution

SLIDE 3

- $Z(\lambda) = \min_{k \in K} (c'x^k + \lambda'(b - Ax^k)), x^k, k \in K$  are extreme points of  $\text{conv}(X)$ .
- $f_k = b - Ax^k$  and  $h_k = c'x^k$ .
- $Z(\lambda) = \min_{k \in K} (h_k + f'_k \lambda)$ , piecewise linear and concave.
- Recall  $\lambda^{t+1} = \lambda^t + \theta_t \nabla Z(\lambda^t)$

### 3.1 Subgradients

SLIDE 4

- Prop:  $f : \Re^n \mapsto \Re$  is concave if and only if for any  $x^* \in \Re^n$ , there exists a vector  $s \in \Re^n$  such that

$$f(x) \leq f(x^*) + s'(x - x^*).$$

- Def:  $f$  concave. A vector  $s$  such that for all  $x \in \Re^n$ :

$$f(x) \leq f(x^*) + s'(x - x^*),$$

is called a **subgradient** of  $f$  at  $x^*$ . The set of all subgradients of  $f$  at  $x^*$  is denoted by  $\partial f(x^*)$  and is called the **subdifferential** of  $f$  at  $x^*$ .

- Prop:  $f : \Re^n \mapsto \Re$  be concave. A vector  $x^*$  maximizes  $f$  over  $\Re^n$  if and only if  $0 \in \partial f(x^*)$ .

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$$\begin{aligned} Z(\lambda) &= \min_{k \in K} (h_k + f'_k \lambda), \\ E(\lambda) &= \{k \in K \mid Z(\lambda) = h_k + f'_k \lambda\}. \end{aligned}$$

Then, for every  $\lambda^* \geq 0$  the following relations hold:

- For every  $k \in E(\lambda^*)$ ,  $f_k$  is a subgradient of the function  $Z(\cdot)$  at  $\lambda^*$ .
- $\partial Z(\lambda^*) = \text{conv}(\{f_k \mid k \in E(\lambda^*)\})$ , i.e., a vector  $s$  is a subgradient of the function  $Z(\cdot)$  at  $\lambda^*$  if and only if  $Z(\lambda^*)$  is a convex combination of the vectors  $f_k$ ,  $k \in E(\lambda^*)$ .

## 3.2 The subgradient algorithm

SLIDE 6

**Input:** A nondifferentiable concave function  $Z(\lambda)$ .

**Output:** A maximizer of  $Z(\lambda)$  subject to  $\lambda \geq 0$ .

**Algorithm:**

- Choose a starting point  $\lambda^1 \geq 0$ ; let  $t = 1$ .
- Given  $\lambda^t$ , check whether  $0 \in \partial Z(\lambda^t)$ . If so, then  $\lambda^t$  is optimal and the algorithm terminates. Else, choose a subgradient  $s^t$  of the function  $Z(\lambda^t)$ .
- Let  $\lambda_j^{t+1} = \max \{\lambda_j^t + \theta_t s_j^t, 0\}$ , where  $\theta_t$  is a positive stepsize parameter. Increment  $t$  and go to Step 2.

### 3.2.1 Step length

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- $\sum_{t=1}^{\infty} \theta_t = \infty$ , and  $\lim_{t \rightarrow \infty} \theta_t = 0$ .
- Example:  $\theta_t = 1/t$ .
- Example:  $\theta_t = \theta_0 \alpha^t$ ,  $t = 1, 2, \dots$ ,  $0 < \alpha < 1$ .
- $\theta_t = f \frac{\hat{Z}_D - Z(\lambda^t)}{\|s^t\|^2}$ , where  $f$  satisfies  $0 < f < 2$ , and  $\hat{Z}_D$  is an estimate of the optimal value  $Z_D$ .
- The stopping criterion  $0 \in \partial Z(\lambda^t)$  is rarely met. Typically, the algorithm is stopped after a fixed number of iterations.

## 3.3 Example

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- $Z(\lambda) = \min \{3 - 2\lambda, 6 - 3\lambda, 2 - \lambda, 5 - 2\lambda, -2 + \lambda, 1, 4 - \lambda, \lambda, 3\}$ ,
- $\theta_t = 0.8^t$ .

	$\lambda^t$	$s^t$	$Z(\lambda^t)$
1.5.00	-3	-9.00	
2.2.60	-2	-2.20	
3.1.32	-1	-0.68	
4.1.83	2	-0.66	
5.1.01	1	-0.99	
6.1.34	1	-0.66	
7.1.60	1	-0.40	
8.1.81	-2	-0.62	
9.1.48	1	-0.52	
10.1.61	1	-0.39	

## 4 Nonlinear problems

SLIDE 9

- $$\begin{aligned} Z_P = \min & f(x) \\ \text{s.t. } & g(x) \leq \mathbf{0}, \\ & x \in X. \end{aligned}$$

- $Z(\lambda) = \min_{x \in X} \{f(x) + \lambda' g(x)\}.$
- $Z_D = \max_{\lambda \geq \mathbf{0}} Z(\lambda).$
- $Y = \{(y, z) \mid y \geq f(x), z \geq g(x), \text{ for all } x \in X\}.$
- $$\begin{aligned} Z_P = \min & y \\ \text{s.t. } & (y, \mathbf{0}) \in Y. \end{aligned}$$
- $Z(\lambda) \leq f(x) + \lambda' g(x) \leq y + \lambda' z, \quad \forall (y, z) \in Y.$
- Geometrically, the hyperplane  $Z(\lambda) = y + \lambda' z$  lies below the set  $Y$ .
- Theorem:

$$\begin{aligned} Z_D = \min & y \\ \text{s.t. } & (y, \mathbf{0}) \in \text{conv}(Y). \end{aligned}$$

### 4.1 Figure

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### 4.2 Example again

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$$X = \{(1, 0)', (2, 0)', (1, 1)', (2, 1)', (0, 2)', (1, 2)', (2, 2)', (1, 3)', (2, 3)'\}.$$

### 4.3 Subgradient algorithm

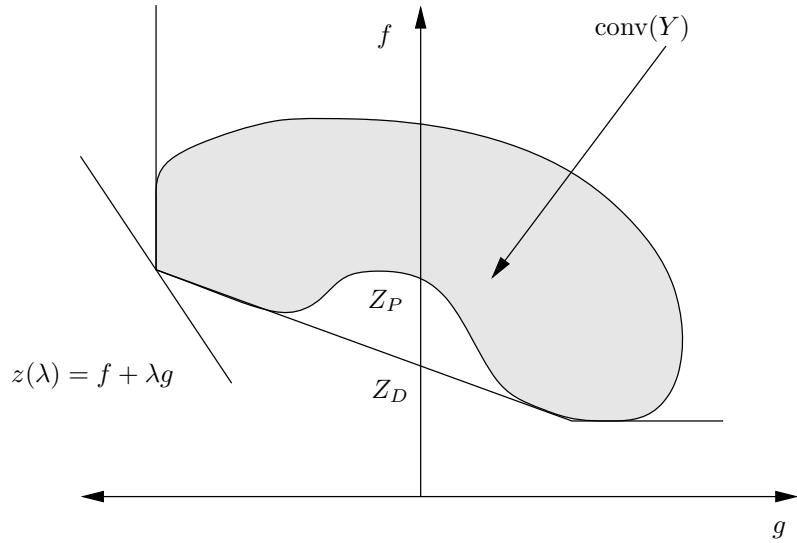
SLIDE 12

**Input:** Convex functions  $f(x), g_1(x), \dots, g_m(x)$  and a convex set  $X$ .

**Output:** An approximate minimizer.

**Algorithm:**

- (Initialization) Select a vector  $\bar{\lambda}$  and solve  $\min_{x \in X} \{f(x) + \lambda' g(x)\}$  to obtain the optimal value  $\bar{Z}$  and an optimal solution  $\bar{x}$ . Set  $x^0 = \bar{x}$ ;  $Z^0 = \bar{Z}$ ;  $t = 1$ .
- (Stopping criterion) If  $(|f(\bar{x}) - \bar{Z}|/\bar{Z}) < \epsilon_1$  and  $(\sum_{i=1}^m |\bar{\lambda}_i|/m) < \epsilon_2$  stop;  
Output  $\bar{x}$  and  $\bar{Z}$  as the solution to the Lagrangean dual problem.



3. **(Subgradient computation)** Compute a subgradient  $s^t$ ;  $\lambda_j^t = \max\{\bar{\lambda}_j + \theta_t s_j^t, 0\}$ , where

$$\theta_t = g \frac{\hat{Z} - Z_{\text{LP}}(\bar{\lambda})}{\|s^t\|^2}$$

with  $\hat{Z}$  an upper bound on  $Z_D$ , and  $0 < g < 2$ . With  $\lambda = \lambda^t$  solve  $\min_{x \in X} \{f(x) + \lambda' g(x)\}$  to obtain the optimal value  $Z^t$  and an optimal solution  $x^t$ .

4. **(Solution update)** Update

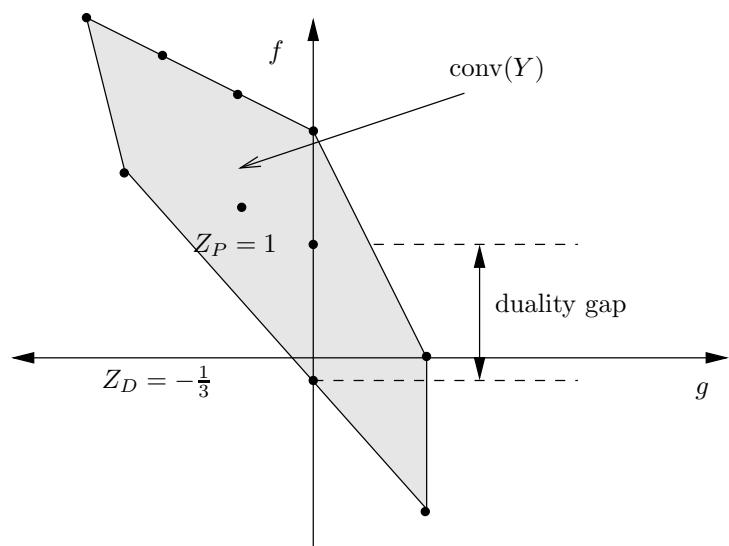
$$\bar{x} \leftarrow \alpha x^t + (1 - \alpha) \bar{x}$$

where  $0 < \alpha < 1$ .

5. **(Improving step)** If  $Z^t > \bar{Z}$ , then

$$\bar{\lambda} \leftarrow \lambda^t, \quad \bar{Z} \leftarrow Z^t;$$

Let  $t \leftarrow t + 1$  and go to Step 2.



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