

15.053/8

April 25, 2013

The Traveling Salesman Problem and Heuristics

Quotes of the day

“Problem solving is hunting. It is savage pleasure and we are born to it.”

-- Thomas Harris

“An algorithm must be seen to be believed.”

-- Donald Knuth

Heuristics

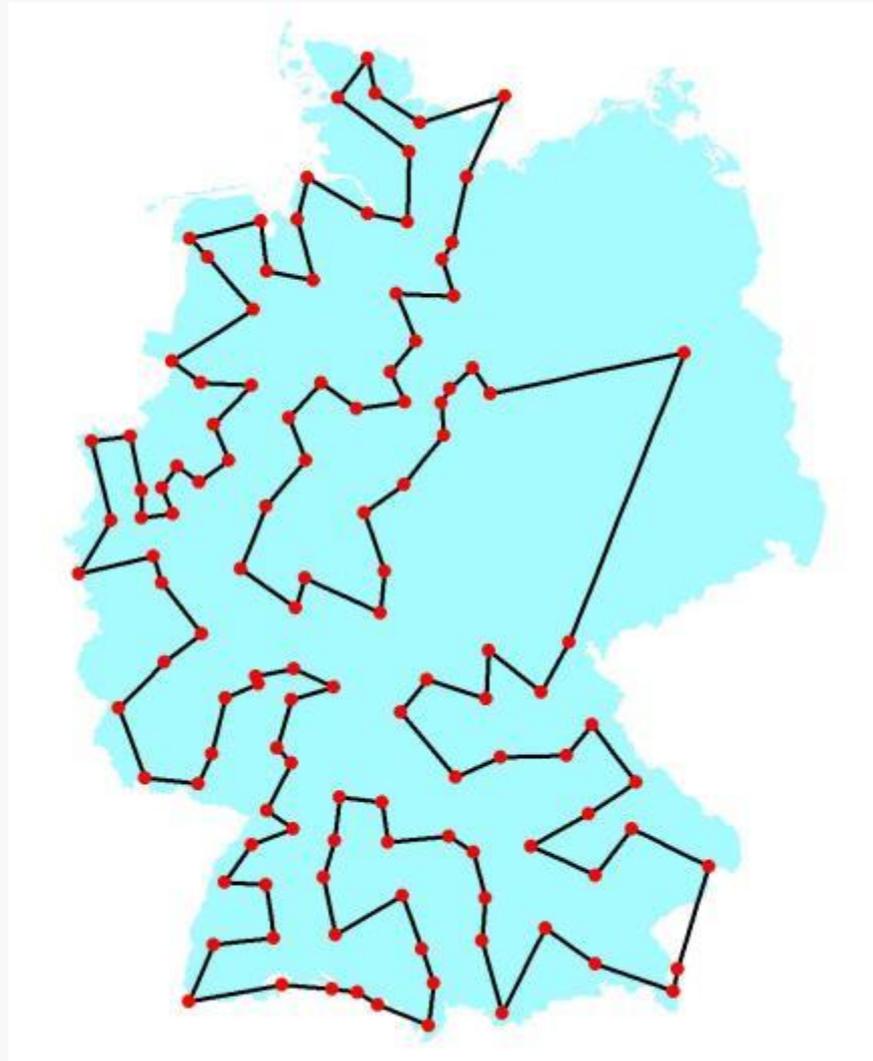
A **heuristic** is a technique designed for solving a problem more quickly when classic methods are too slow (from Wikipedia).

Today's lecture:

- Heuristics illustrated on the traveling salesman problem.
- Design principles for heuristics
- Chances for practice

Traveling Salesperson Problem (TSP)

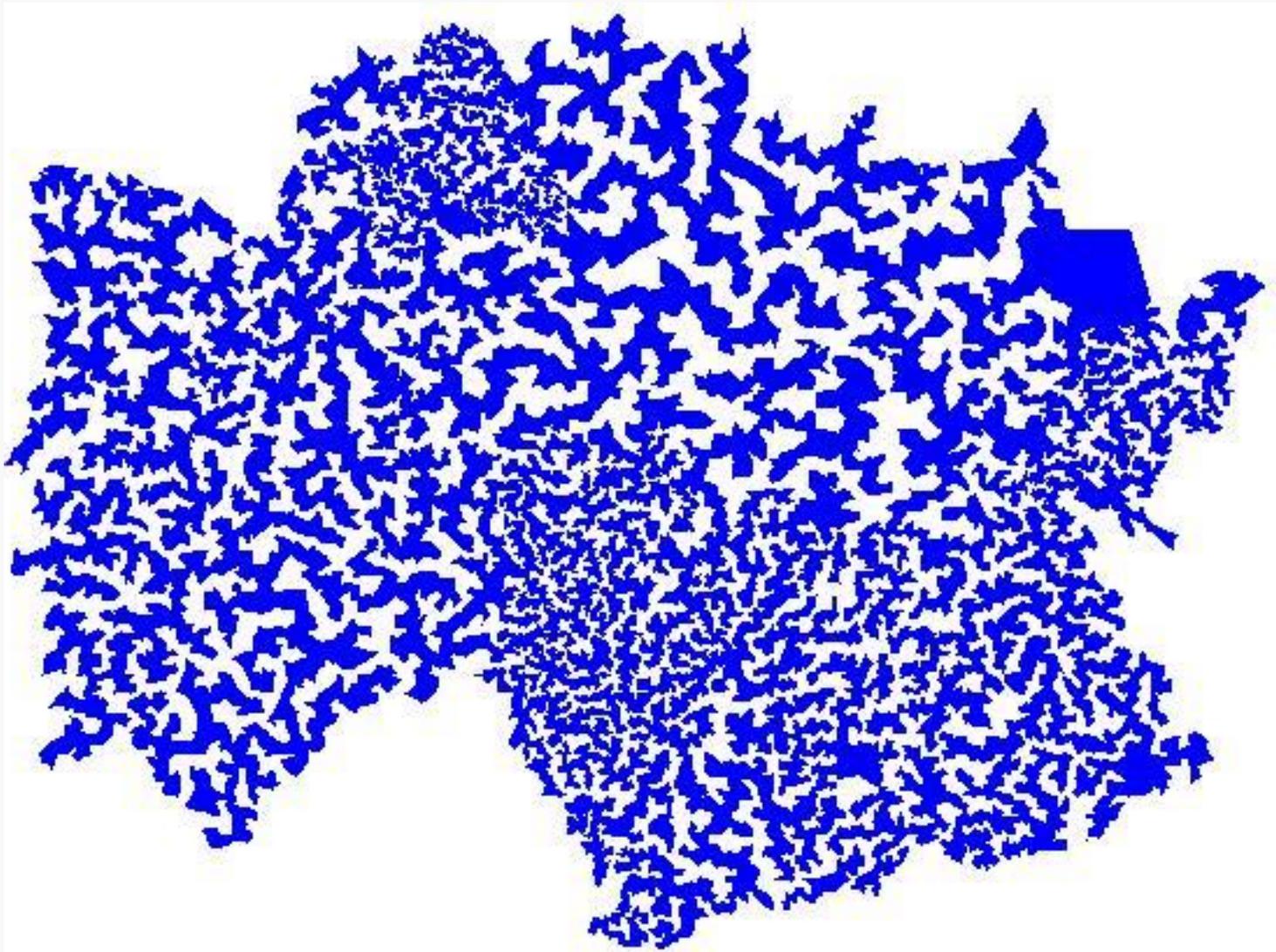
- Find the shortest distance tour passing through each node of the network exactly once.
- c_{ij} = distance from i to j .



<http://www.math.uwaterloo.ca/tsp/>

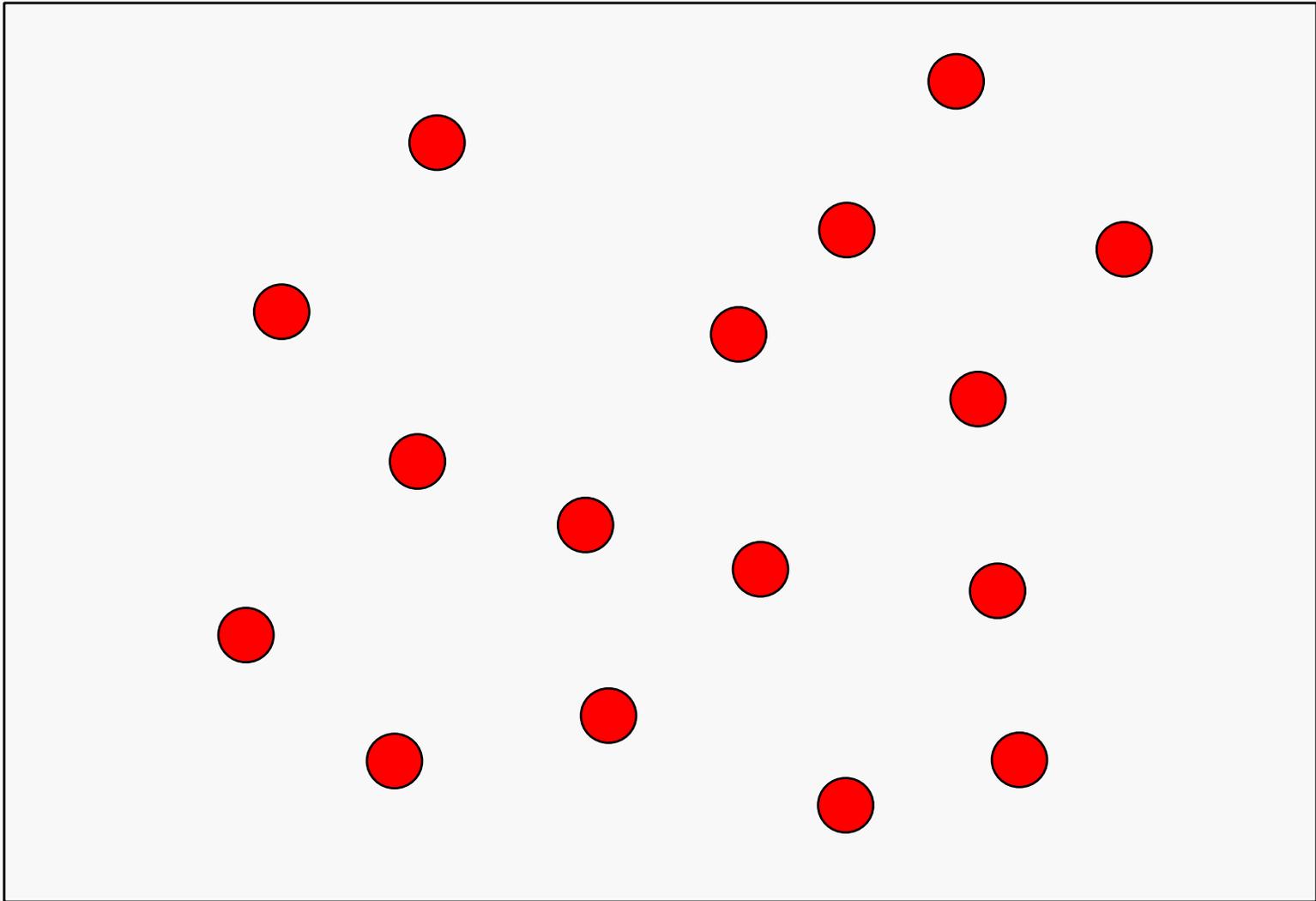
Courtesy of William Cook. Used with permission.

15,112 City Optimal Tour in Germany (rotated)



Courtesy of William Cook. Used with permission.

Exercise: Try to find the best tour.



On solving hard problems

- **How did you select your tour?**
 - it relied on visualization
 - perhaps you took a holistic view
- **How can we develop computer heuristics for solving hard problems?**

The TSP is a hard problem

- NP-hard (also NP-complete)

Have you seen NP-hardness or NP-completeness before?

- 1. Yes.**
- 2. No.**

The TSP is a hard problem

- **There is no known polynomial time algorithm. Cannot bound the running time as less than n^k for any fixed integer k (say $k = 15$).**
- **If there were a polynomial time algorithm, there would be a polynomial time algorithm for every NP-complete problem.**
- **Question: what does one do with a hard problem?**

100 n¹⁵ vs. n!

Suppose that we could carry out 1 sextillion steps per second (10²¹).

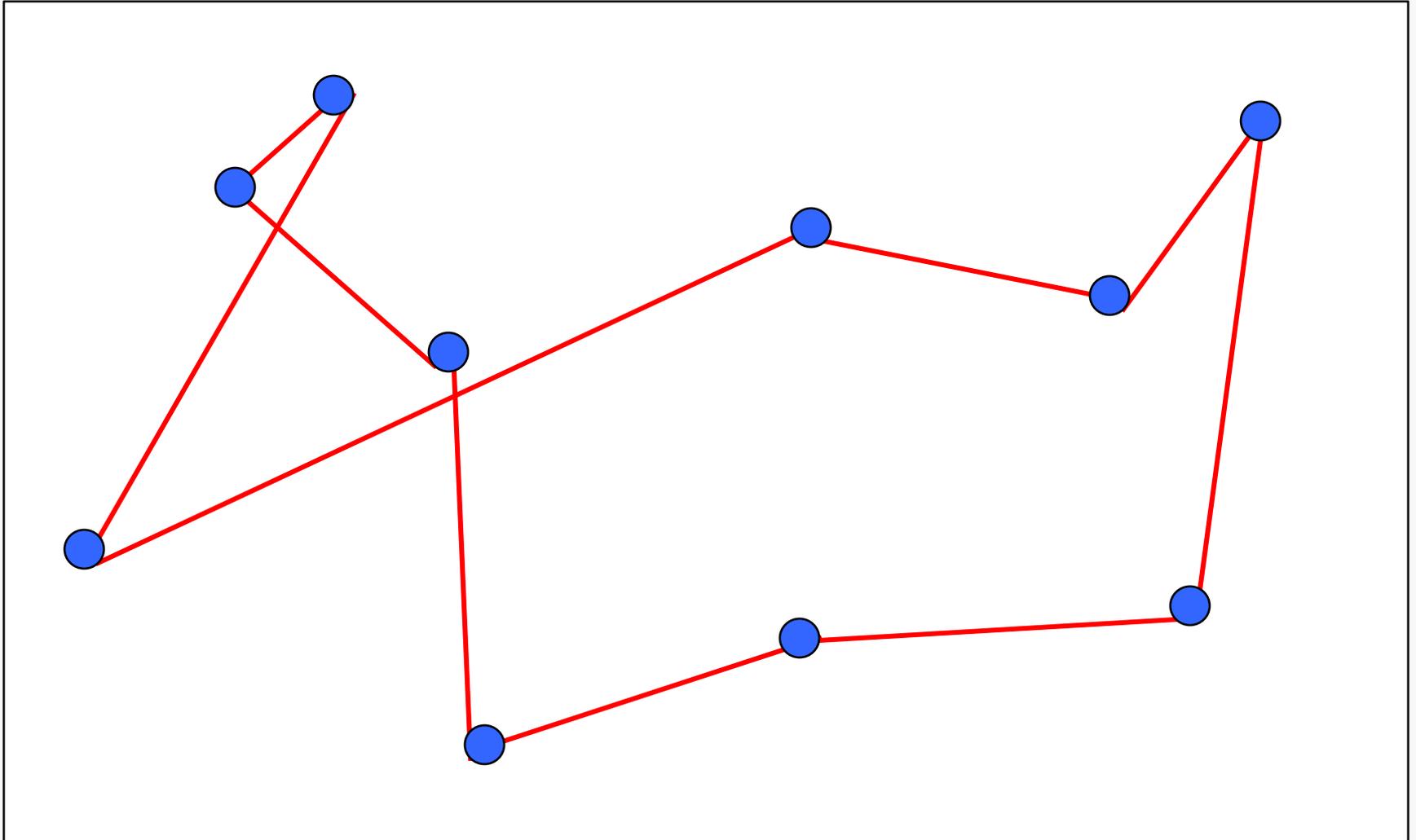
- n = number of cities. (n-1)! tours.
- compare time for 100 n¹⁵ steps vs. n! steps.
 - Remark: 100 n¹⁵ steps is NOT practically efficient.

<u># of cities</u>	<u>100 n¹⁵ steps</u>	<u>n! steps</u>
n = 20,	3.27 seconds	2.4 milliseconds
n = 25,	1.5 minutes	4.2 hours
n = 30	24 minutes	8,400 years
n = 35	4 hours	326 billion years

Two types of Heuristics

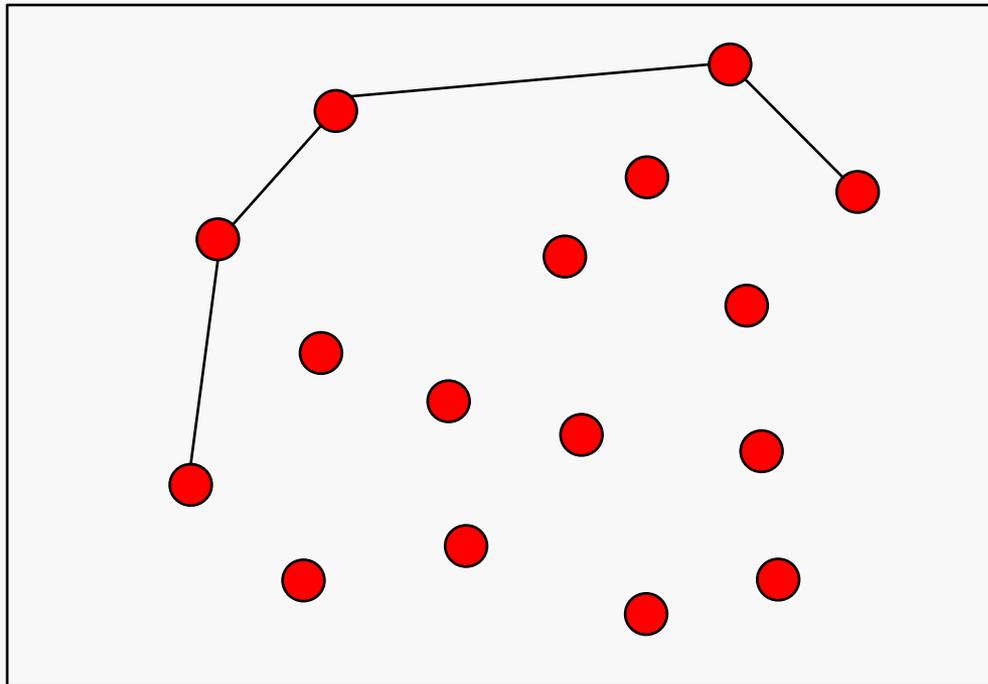
- **Construction heuristics: builds a solution from scratch (starting with nothing).**
 - **Often called “greedy heuristics”. Each step looks good, but it doesn’t look ahead.**
- **Improvement heuristics (neighborhood search): starts with a solution, and then tries to improve the solution, usually by making small changes in the current solution.**

An easy construction heuristic: Nearest unvisited neighbor



Can we do better in construction?

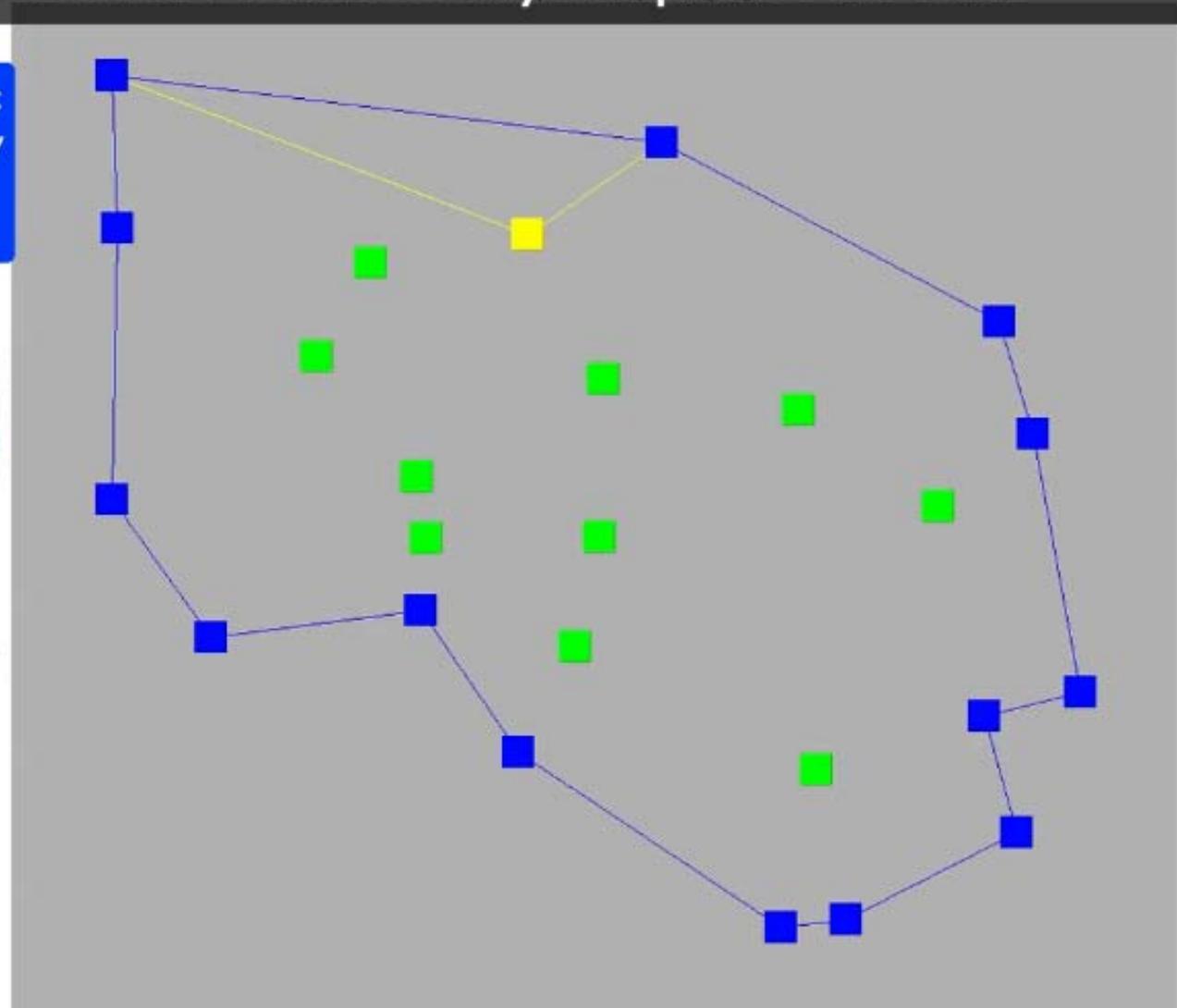
- **Class exercise:** try to develop a heuristic in which we add one city at a time, but the next city can be added anywhere in the tour (not just the beginning or the end.)
 - Below is the beginning part of a tour



TSP Insertion Heuristic by Stephan Mertens

Cheapest Insertion Heuristic
This heuristic selects the city whose addition to the tour has least cost.

- new
- reset
- step
- run
- Nodes:
 - 10
 - 15
 - 20
 - 25
- solve



- Speed
 - fast
 - med.
 - slow
- Insert
 - cheapest
 - farthest

Courtesy of Stephan Mertens. Used with permission.

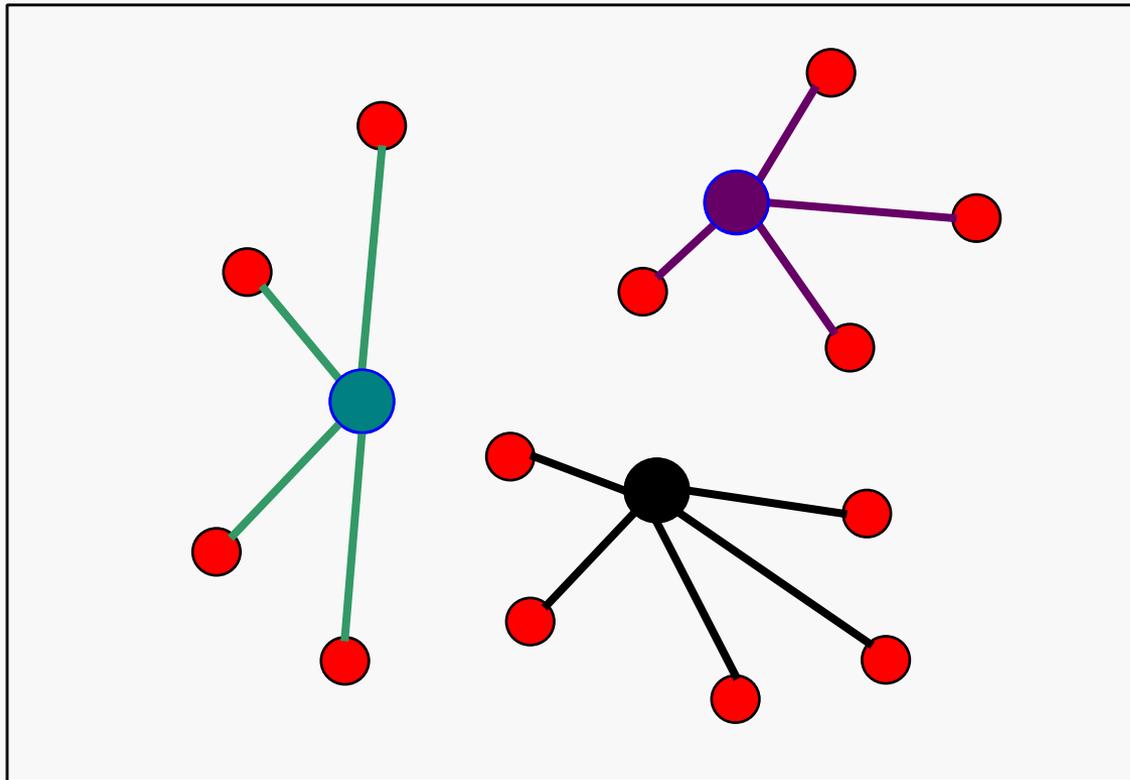
Which do you believe will give shorter length tours:

1. The nearest neighbor heuristic
2. An insertion based heuristic

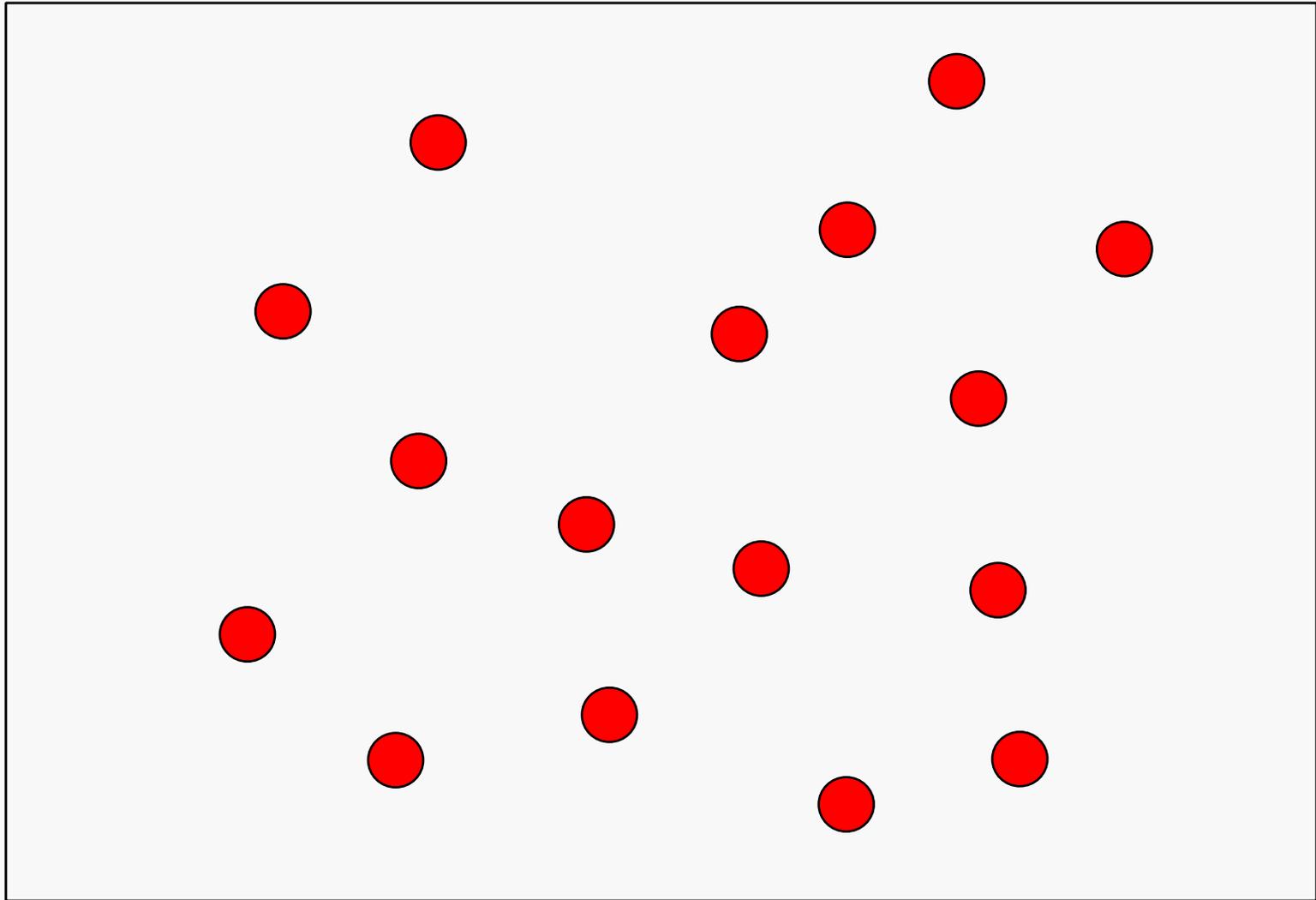
Facility location problems.

Choose K facilities so as to minimize total distance from customers to their closest facility.

- example with three facilities



Exercise: try developing a good solution where there are 2 facilities



Exercise: Develop a construction heuristic for the facility location problem

- **Data:** locations in a city.
- c_{ij} = distance from i to j

Mental Break

Improvement heuristics

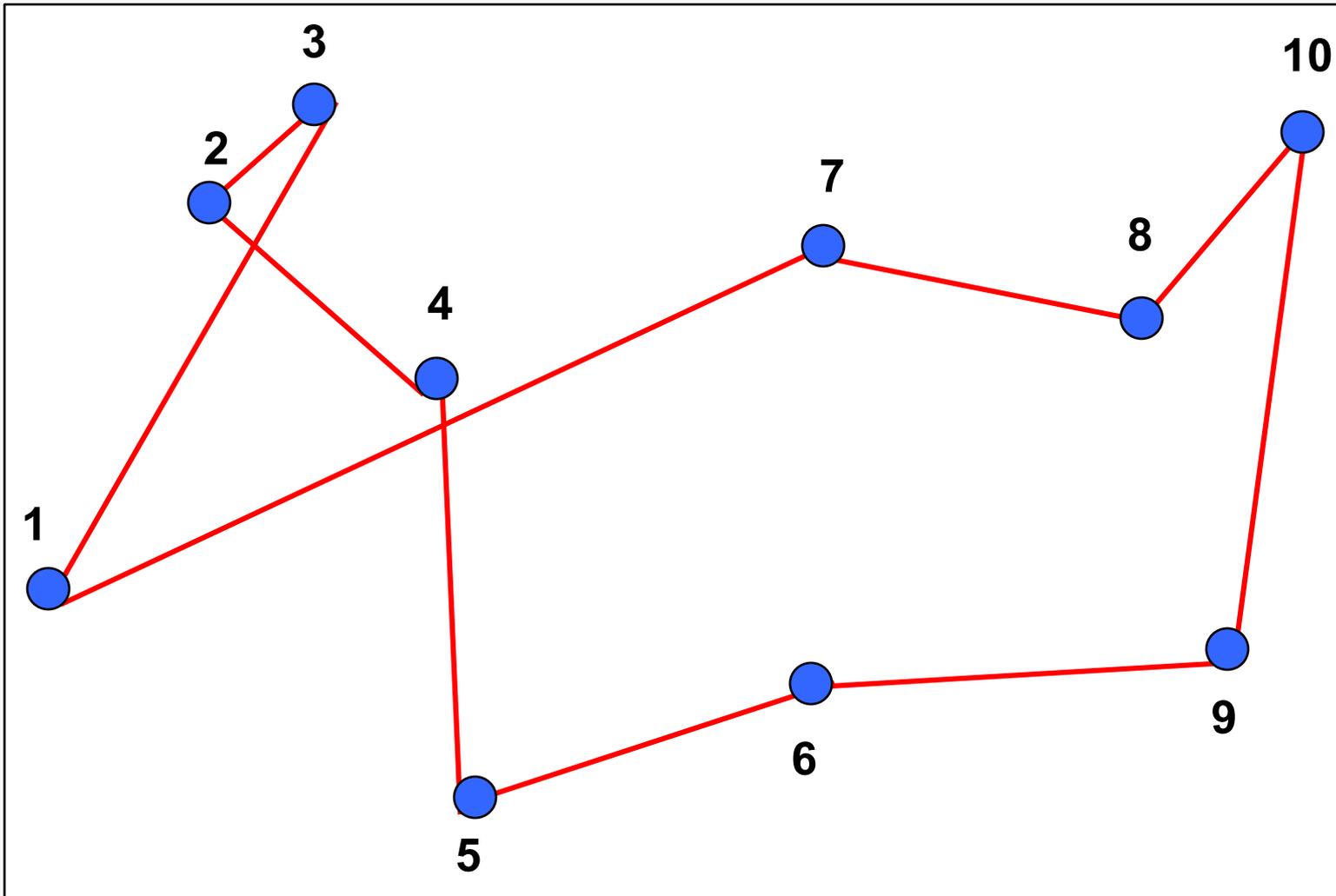
- Improvement heuristics start with a feasible solution and look for an improved solution that can be found by making a very small number of changes.
 - This will be made more formal
- Two TSP tours are called 2-adjacent if one can be obtained from the other by deleting two edges and adding two edges.

2-opt neighborhood search

- A TSP tour T is called **2-optimal** if there is no 2-adjacent tour to T with lower cost than T .
- **2-opt heuristic.** Look for a 2-adjacent tour with lower cost than the current tour. If one is found, then it replaces the current tour. This continues until there is a 2-optimal tour.

An improvement heuristic: 2 exchanges.

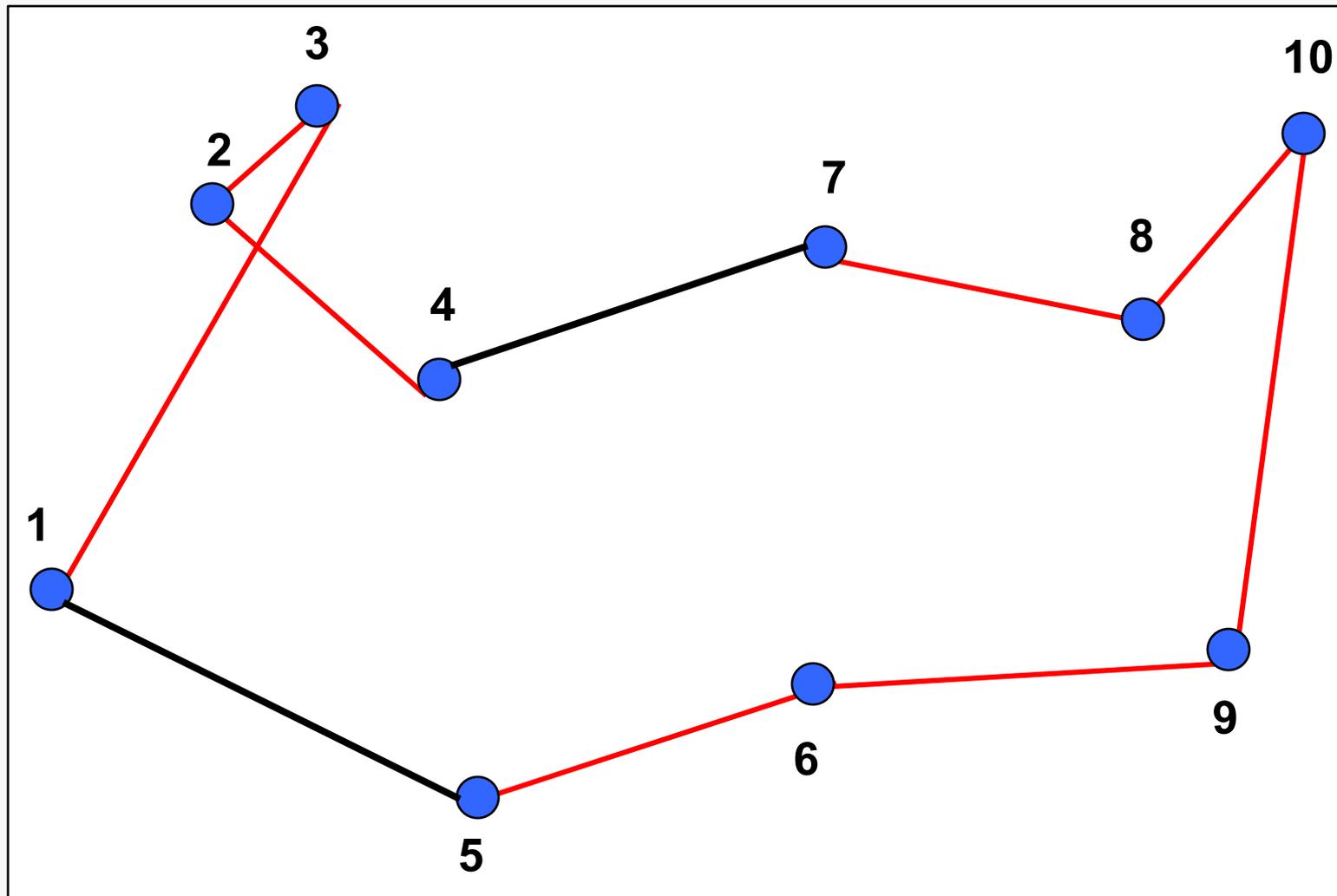
Look for an improvement obtained by deleting two edges and adding two edges.



T
1
3
2
4
5
6
9
10
8
7

After the two exchange

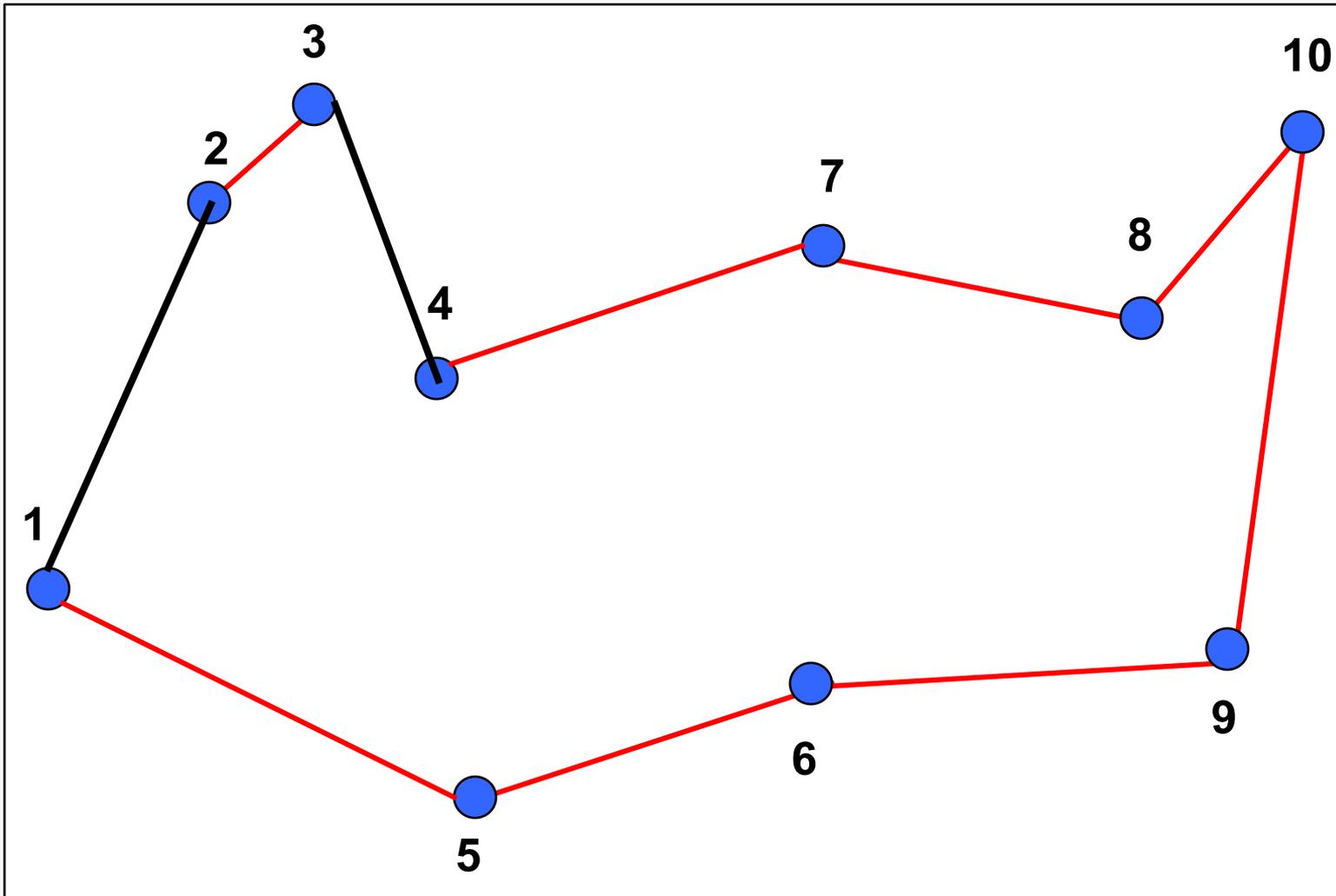
Deleting arcs (4,7) and (5, 1) flips the subpath from node 7 to node 5.



T_1	T_2
1	1
3	3
2	2
4	4
5	7
6	8
9	10
10	9
8	6
7	5

After the two exchange

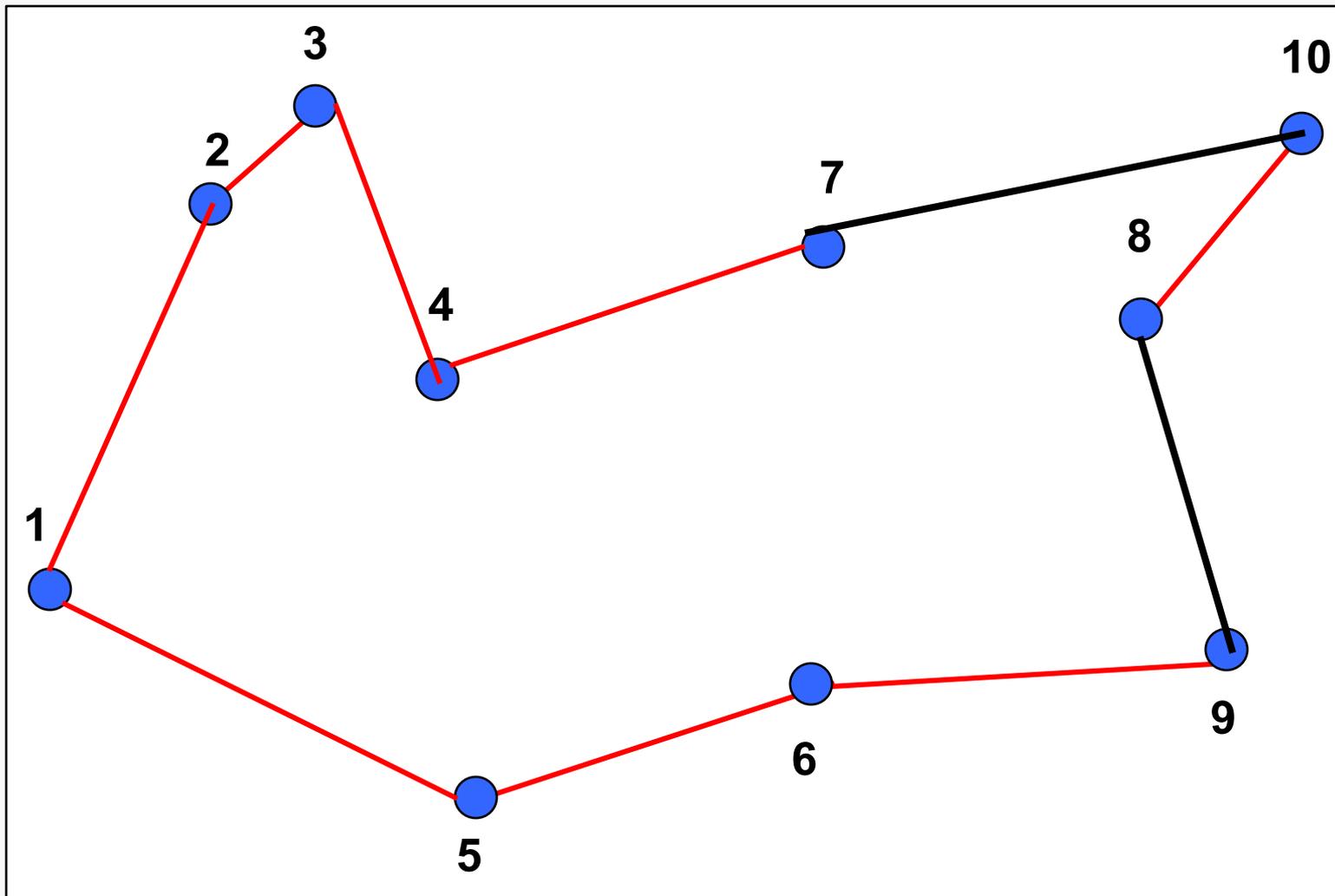
Deleting arcs (1,3) and (2, 4) flips the subpath from 3 to 2.



T_2	T_3
1	1
3	2
2	3
4	4
7	7
8	8
10	10
9	9
6	6
5	5

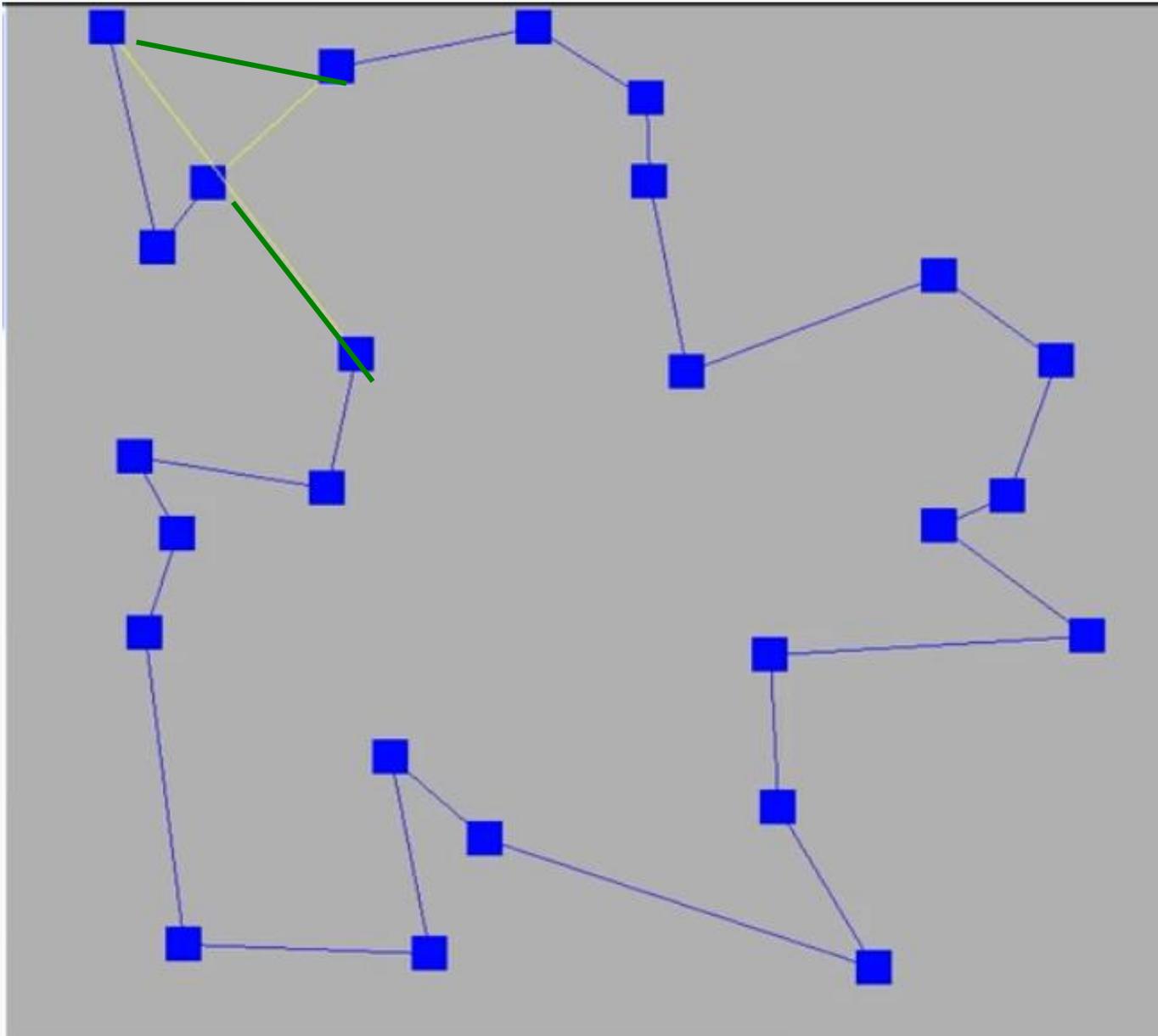
After the final improving 2-exchange

Deleting arcs (7,8) and (10, 9) flips the subpath from 8 to 10.



T_3	T_4
1	1
2	2
3	3
4	4
7	7
8	10
10	8
9	9
6	6
5	5

2-exchange heuristic (also called 2-opt)



3-opt neighborhood

- Two TSP tours are called **3-adjacent** if one can be obtained from the other by deleting three edges and adding three edges.
- A TSP tour T is called **3-optimal** if there is no 3-adjacent tour to T with lower cost than T .
- **3-opt heuristic.** Look for a 3-adjacent tour with lower cost than the current tour. If one is found, then it replaces the current tour. This continues until there is a 3-optimal tour.

On Improvement Heuristics

Improvement heuristics are based on searching a “neighborhood .“ Let $N(T)$ be the neighborhood of tour T .

In this case, the $N(T)$ consists of all tours that can be obtained from T deleting two arcs and inserting two arcs.

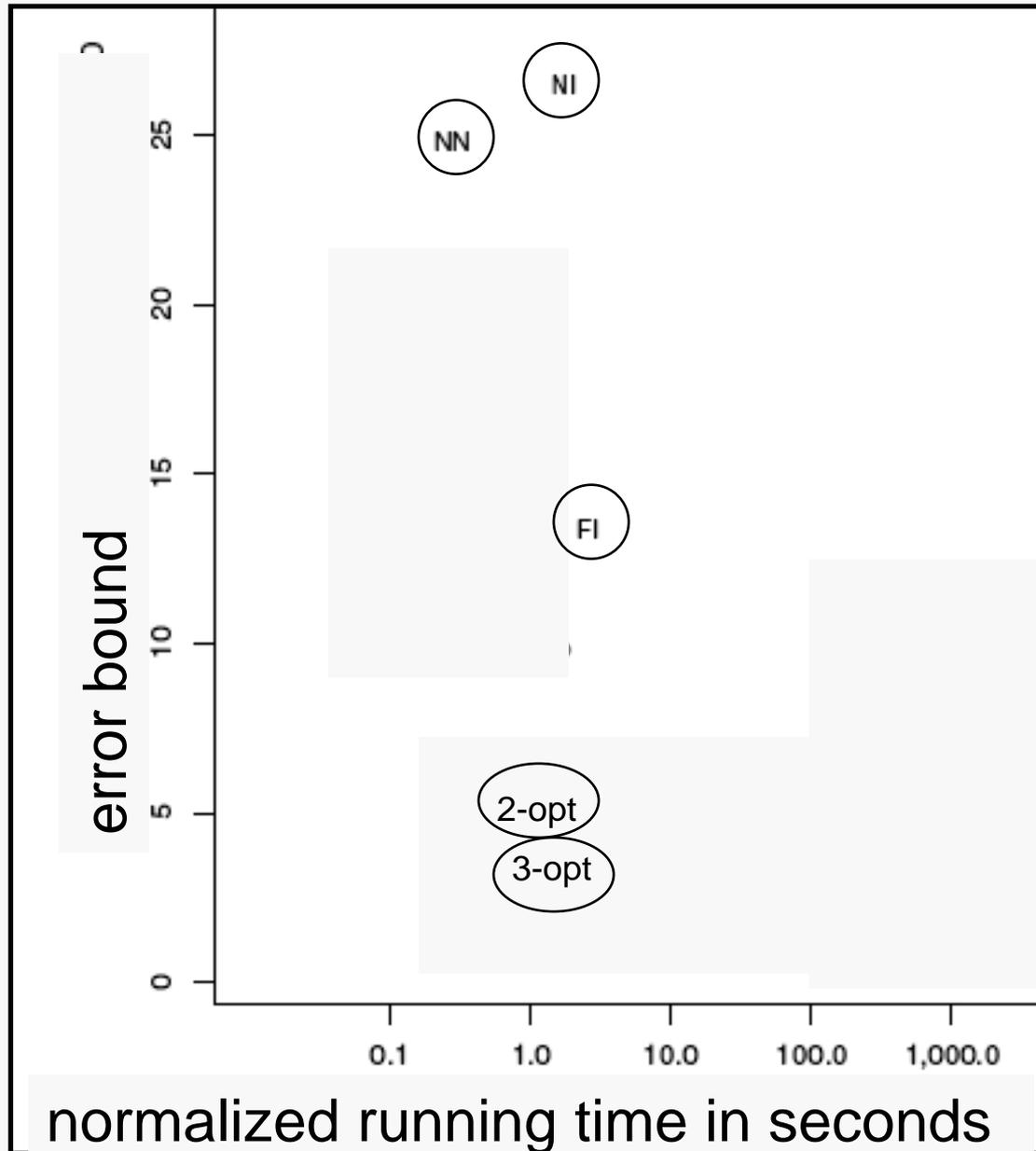
Improvement heuristic:

start with tour T

if there is a tour $T' \in N(T)$ with $c(T') < c(T)$, then
replace T by T' and repeat

otherwise, quit with a *locally optimal solution*.

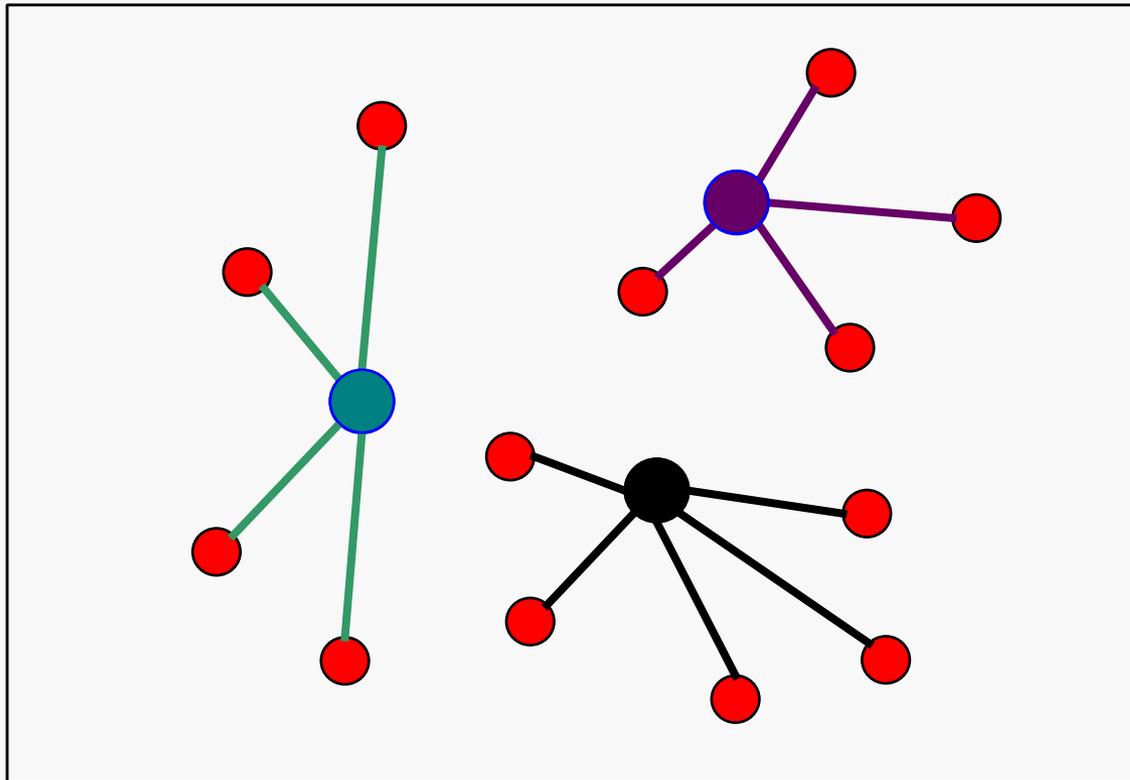
How good are improvement heuristics?



Implementers had to be very clever to achieve these running times.

Facility location problems.

Class exercise. Suppose we want to solve a facility location problem with 3 facilities. Design a neighborhood search heuristic.



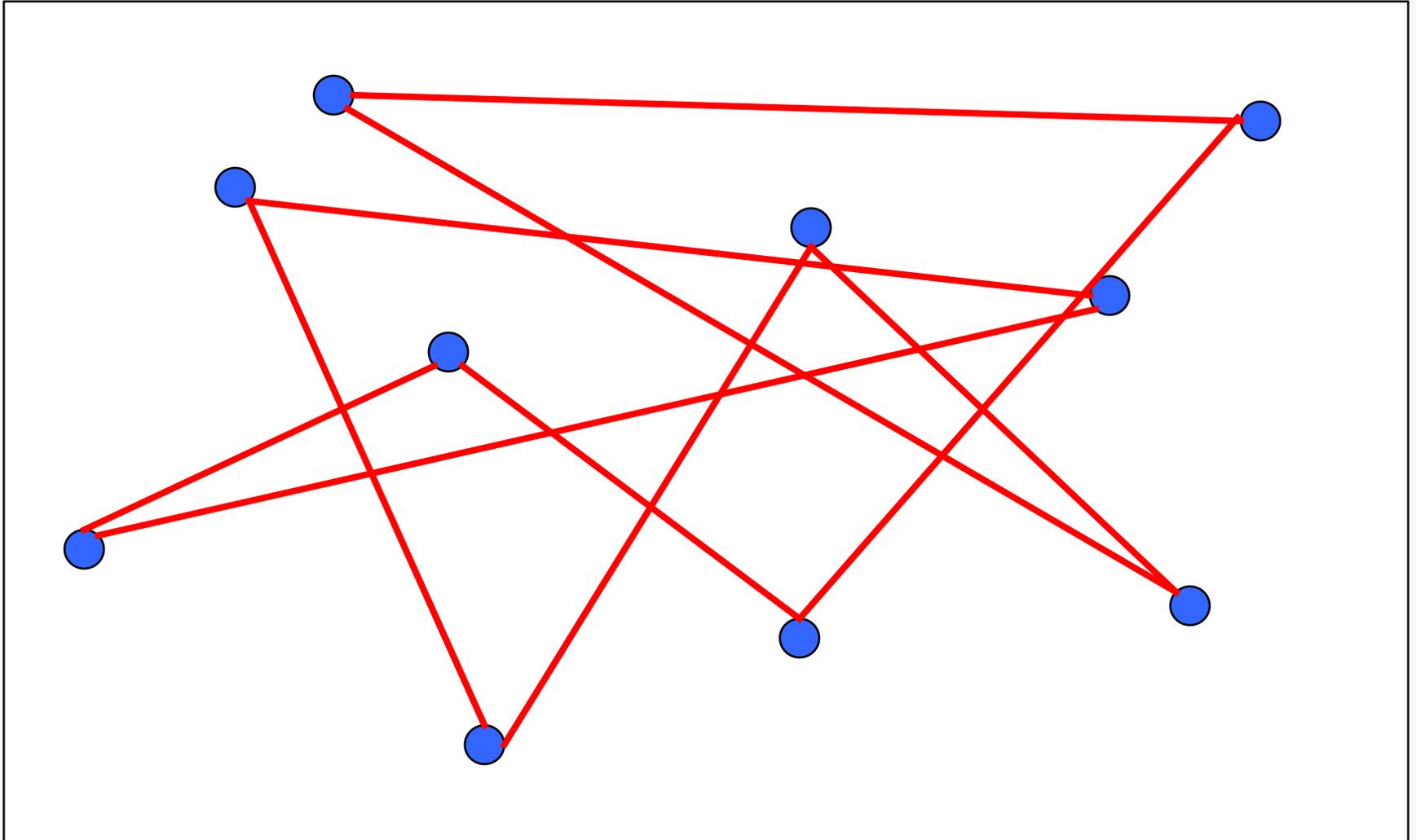
Using Randomization

- **An important idea in algorithm development: randomization**
- **Randomization in neighborhood improvement: a way of using the same approach multiple times and getting different answers. (Then choose the best).**
- **Simulated Annealing: randomization in order to have an approach that is more likely to converge to a good solution**

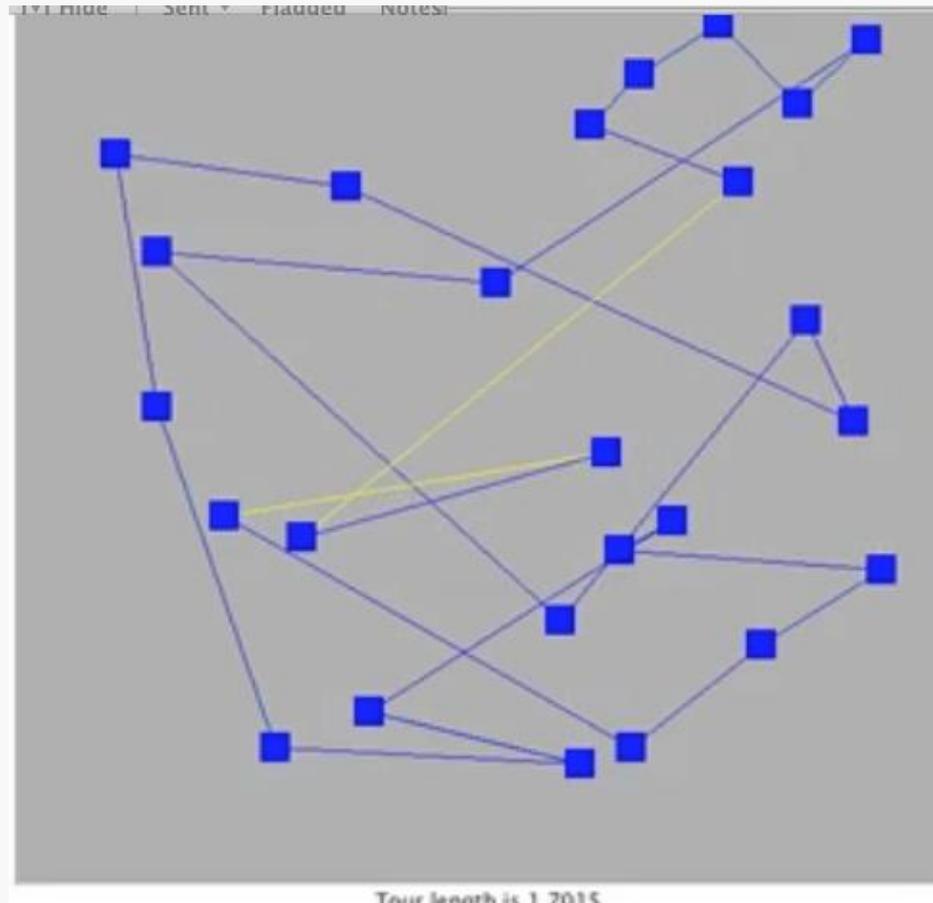
On the use of randomization

- **Remark: 2-exchanges will behave differently depending on the starting solution.**
- **Randomization based heuristic:**
- **Start with a random tour**
- **use the 2-exchange neighborhood until obtaining a local optimum.**

One difficulty: random tours are terrible.



2-opt heuristic starting from a random solution.



Another use of randomization

- **Replace the nearest neighbor tour with the following: at each iteration, visit either the closest neighbor or the second or third closest neighbors. Choose each with 1/3 probability.**
- **This generates a random tour that is “pretty good” and may be a better starting point than a totally random tour.**

Other approaches to heuristics

- **The metaphor based approach to the design of heuristics**
 - **simulated annealing**
 - **genetic algorithms**
 - **neural networks**
 - **ant swarming**
- **That is, look for something that seems to work well in nature, and then try to simplify it so that it is practical and helps solve optimization problems.**

Simulated Annealing

A randomization heuristic based on neighborhood search that permits moves that make a solution worse. It is based on an analogy with physical annealing.

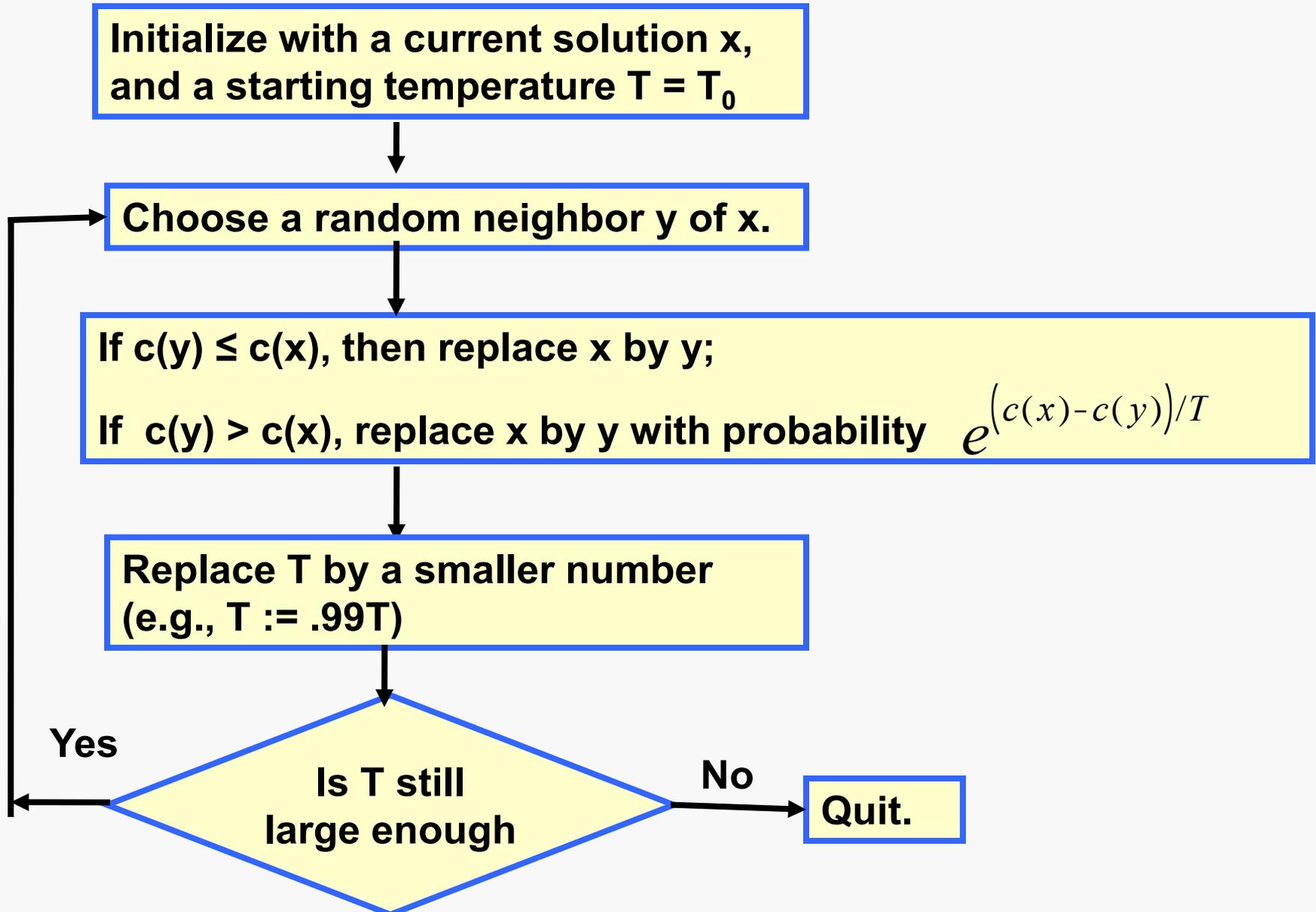
Photo removed due to
copyright restrictions.

**To take a hot material
and have it reach a low
energy state, one
should cool it slowly.**

A glass annealing oven.
www.carbolite.com

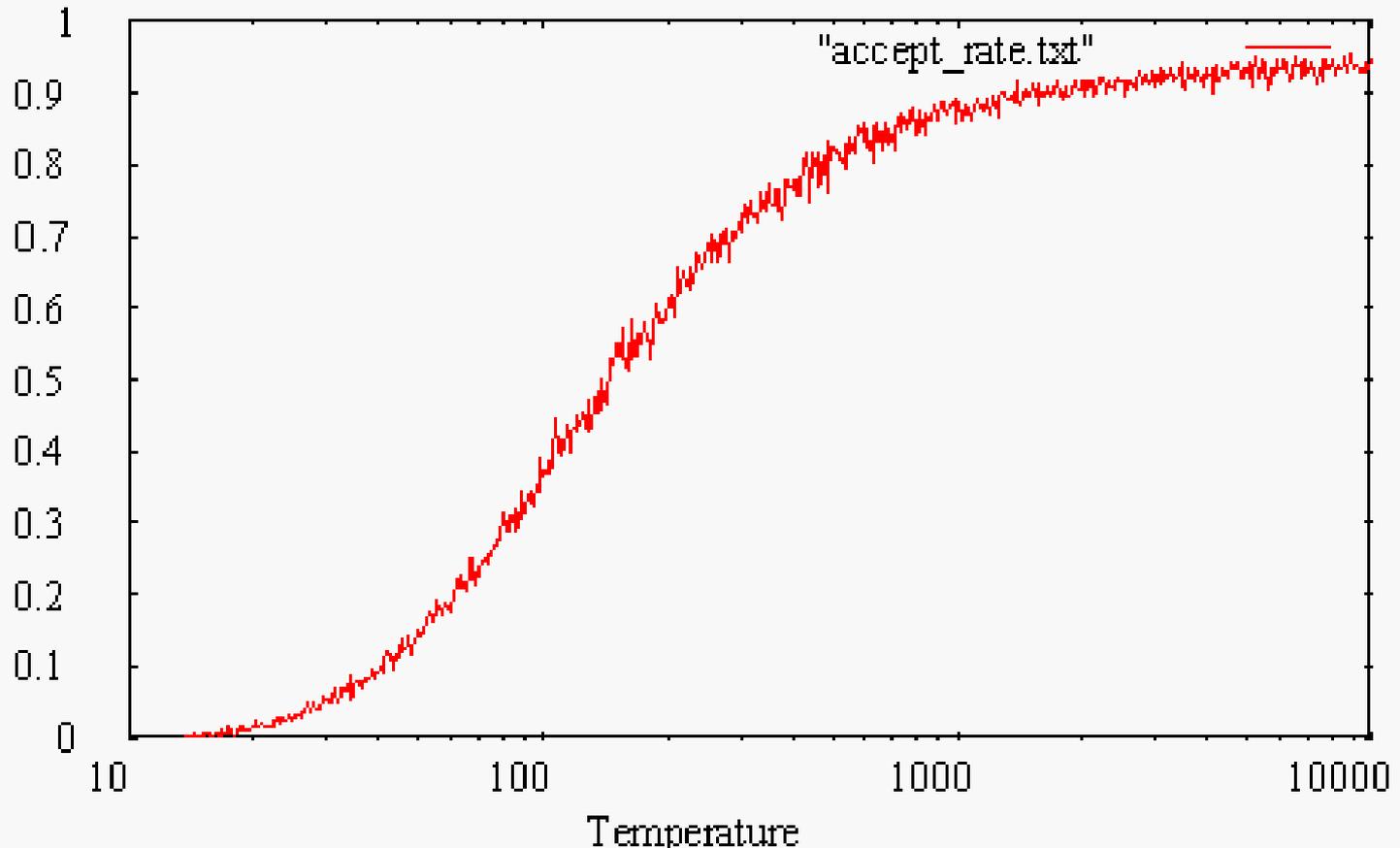
See Eglese, R. W. "Simulated Annealing: A tool for Operational Research."
European Journal of Operational Research 46 (1990) 271-281. ([PDF](#))

Simulated Annealing: a variation on local search.



A typical acceptance rate as a function of temp

Accept Rate @each Temperature

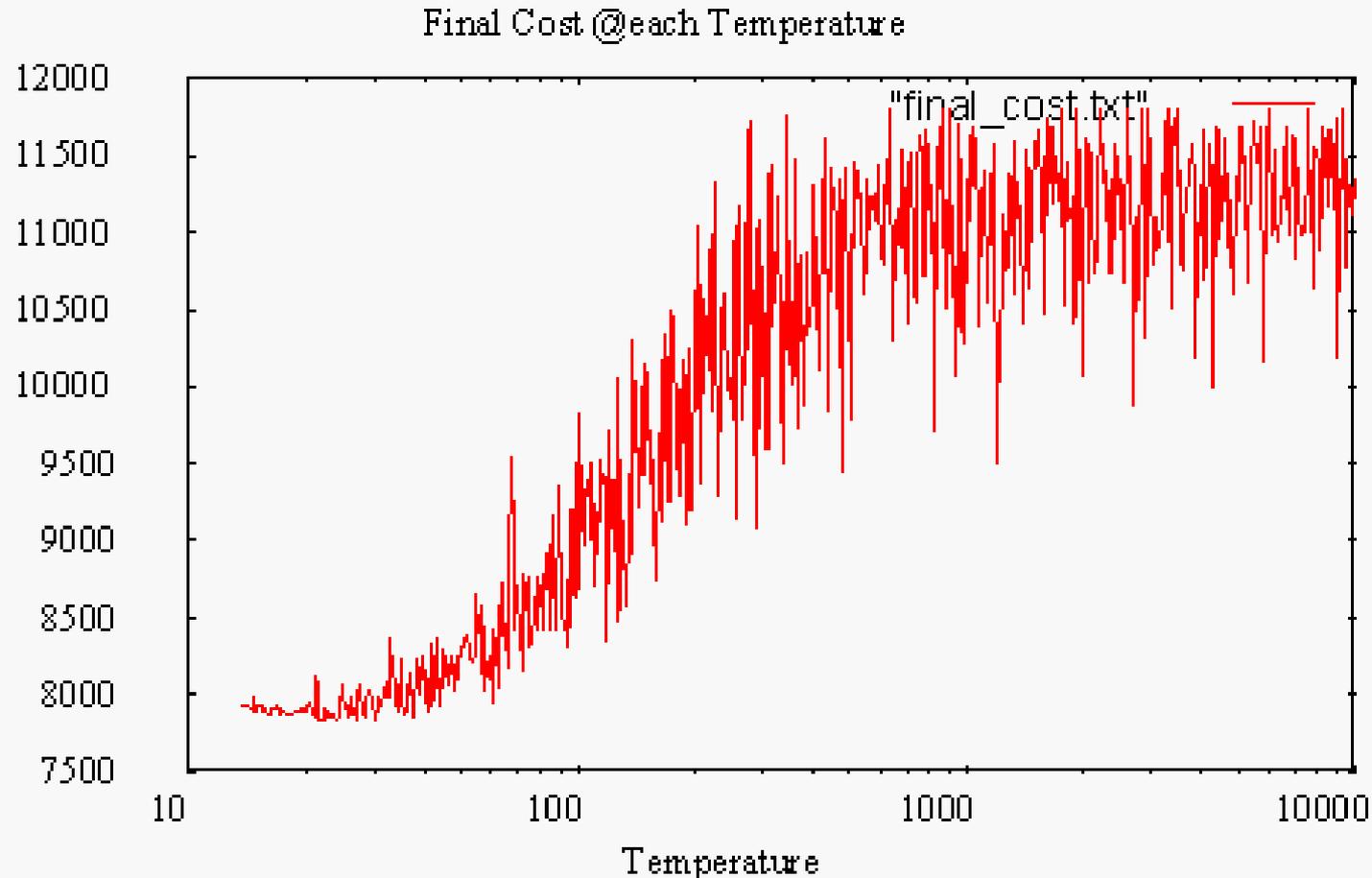


A network design problem for wireless networks.

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A typical graph of Temp vs. cost.



A network design problem for wireless networks.

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Summary

- **Two types of heuristics**
- **Use of randomization**
- **Simulated annealing**
- **Goal of this lecture: let you get started in solving hard problems.**

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