

# Problem Set 7 Solution

17.881/882

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## 1 Gibbons 2.3 (p.131)

Let us consider the three-period game first.

### 1.1 The Three-Period Game

The structure of the game as was described in section 2.1D (pp.68-71). Let us solve the game backwards.

#### 1.1.1 Stage (2b)

Player 1 accepts player 2's proposal  $(s_2, 1 - s_2)$  if and only if the following condition is satisfied:

$$s_2 \geq \delta_1 s$$

#### 1.1.2 Stage (2a)

Conditional on player 1 accepting the offer, player 2 maximises his/her payoff by offering  $s_2 = \delta_1 s$ . Then, player 2 gets  $\delta_2(1 - s_2) = \delta_2(1 - \delta_1 s)$ .

Any rejected offer leads 2 to get a payoff of  $\delta_2^2(1 - s)$ . Player 2 is better off with a proposal that is accepted if and only if:

$$\begin{aligned} d_2 &= \delta_2(1 - \delta_1 s) - \delta_2^2(1 - s) \\ &= \delta_2[1 - \delta_2 + s(\delta_2 - \delta_1)] \geq 0 \end{aligned}$$

If  $\delta_2 \geq \delta_1$ , we have  $d_2 \geq \delta_2[1 - \delta_2] > 0$  since  $0 < \delta_2 < 1$ .

If  $\delta_2 < \delta_1$ , we have  $d_2 \geq \delta_2[1 - \delta_1] > 0$  since  $0 < \delta_1, \delta_2 < 1$ .

Either way, we have that player 2 prefers to have his/her proposal accepted, and offers  $s_2 = \delta_1 s$

### 1.1.3 Stage (1b)

Player 2 accepts player 1's proposal  $(s_1, 1 - s_1)$  if and only if:

$$\begin{aligned} 1 - s_1 &\geq \delta_2(1 - s_2) \\ s_1 &\leq 1 - \delta_2[1 - \delta_1 s] \end{aligned}$$

### 1.1.4 Stage (1a)

Conditional on player 2 accepting the offer, player 1 maximises his/her payoff by offering  $s_1 = 1 - \delta_2[1 - \delta_1 s]$ .

Any rejected offer leads 1 to get a payoff of  $\delta_1 s_2 = \delta_1^2 s$ . Player 1 is better off with a proposal that is accepted if and only if:

$$\begin{aligned} d_1 &= 1 - \delta_2[1 - \delta_1 s] - \delta_1^2 s \\ &= 1 - \delta_2 + s\delta_1(\delta_2 - \delta_1) \geq 0 \end{aligned}$$

If  $\delta_2 \geq \delta_1$ , we have  $d_1 \geq 1 - \delta_2 > 0$  since  $0 < \delta_2 < 1$ .

If  $\delta_2 < \delta_1$ , we have  $d \geq (1 - \delta_2 + \delta_1)(1 - \delta_1) > 0$  since  $0 < \delta_1, \delta_2 < 1$ .

Either way, we have that player 1 prefers to have his/her proposal accepted, and offers  $s_1 = 1 - \delta_2[1 - \delta_1 s]$

The outcome of the game is that players 1 and 2 agree on the distribution  $(s_1^*, 1 - s_1^*) = (1 - \delta_2[1 - \delta_1 s], \delta_2[1 - \delta_1 s])$ .

## 1.2 The Infinite-Horizon Game

Let  $s$  be a payoff that player 1 can get in a backwards-induction of the game as a whole, and  $s_H$  the maximum value of  $s$ . Imagine using  $s$  as the third-payoff to Player 1. Player 1's first-period payoff is a function of  $s$ , namely  $f(s) = 1 - \delta_2[1 - \delta_1 s]$ . Since this function is increasing in  $s$ ,  $f(s_H)$  is the highest possible first-period payoff, so  $f(s_H) = s_H$ . Then

$$\begin{aligned} 1 - \delta_2[1 - \delta_1 s_H] &= s_H \\ \iff s_H &= \frac{1 - \delta_2}{1 - \delta_1 \delta_2} \end{aligned}$$

A parallel argument shows that  $f(s_L) = s_L$ , where  $f(s_L)$  is the lowest payoff that player 1 can achieve in any backwards-induction of the game as a whole. Therefore, the only value of  $s$  that satisfies  $f(s) = s$  is  $\frac{1 - \delta_2}{1 - \delta_1 \delta_2}$ . Thus  $s_H = s_L = s^*$ , so there is a unique backwards-induction outcome of the game as a whole: A distribution

$$(s^*, 1 - s^*) = \left\{ \frac{1 - \delta_2}{1 - \delta_1 \delta_2}, \frac{\delta_2(1 - \delta_1)}{1 - \delta_1 \delta_2} \right\}$$