

# Lecture 19: Quantum Games

8.371 p.115  
4/25/06

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- 1) Classical games
- 2) PQ Penny Flipover
- 3) Q. Prisoner's Dilemma
- 4) Tragedy of the Commons
- 5) Q. Public Goods Game

Games field founded by von Neumann, Morgenstern (1950s)

Multiperson decision theory

- Analysis of decision making process, assuming rational behavior (i.e. each person maximizes rewards, profits, incomes, subjective benefits)

Example: Prisoner's Dilemma (Albert Tucker, 1950)

2 burglars caught, separated by police (no communication)  
Each can either cooperate w/ partner (stay silent) or defect (confess)

Rewarded  $r=1$   
Punished  $p=1$   
Tempted  $t=5$   
Sucker's payoff  $s=0$

Payoff:

		Bob		
		C	D	
AQ	C	3,3	0,5	(A,B)
	D	5,0	1,1	

Goal: Maximize payoff

Clear optimum choice: D, D

Def: Dominant strategy earns a player a larger payoff than any other strategy, regardless of what other players do.

Example 2: Price Wars

		Shor's Widgets			(C, S)
		#1	#2	#3	
Chvane's Gadgets	#1	0, 0	50, -10	40, -20	
	#2	-10, 50	20, 20	90, 10	
	#3	-20, 40	10, 90	50, 50	

Trajectory always leads to 0,0

No dominant strategy here!

(Player's choices depend very much on other's)

∃ local max: At #1, #1, neither wants to change

Def: Nash Equilibrium is a set of strategies (one for each player) s.t. no player has incentive to change his/her action (partial deriv = 0)

Example 3: Lefty-Righty, Uppy-Downy

	L	R
U	8, 8	0, 6
D	6, 0	7, 7

2 Nash equil:  $\{LU, RD\}$

(Can move along straight line in table,  
but not diagonally)

Example 4: 3-Player Lefty-Righty

	L	R
U	0, 0, 10	-5, -5, 10
D	-5, -5, 0	1, 1, -5

Matrix A

	L	R
U	-2, -2, 0	-5, -5, 0
D	-5, -5, 0	-1, -1, 5

Matrix B

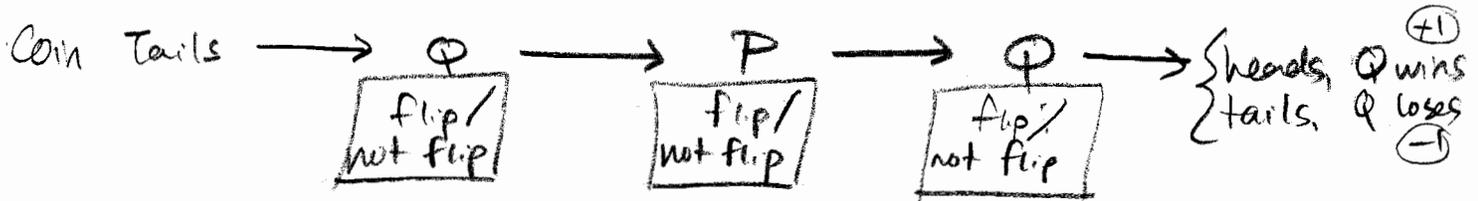
(3 players: Col, Row, Matrix)

Nash equil:  $ULA$

Vulnerable to coalition  $\rightarrow DRB$

Nash equil: Unilaterally optimal

## II PQ Penny Flipover



		$\Phi$			
		NN	NF	FN	FF
$P$	N	-1	1	1	-1
	F	1	-1	-1	1

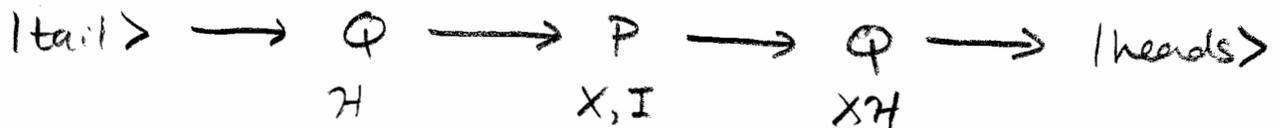
No Nash equil

Best strategy: mixed

$$P = \frac{1}{2} N, \frac{1}{2} F$$

$$Q = \frac{1}{4} \{ NN, FN, NF, FF \}$$

### Quantum Version:

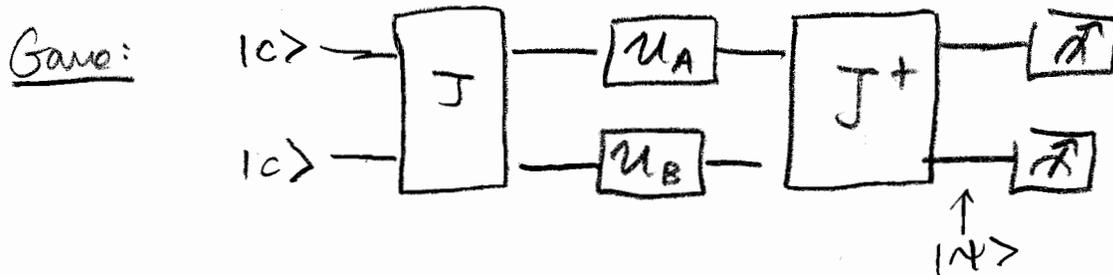


Q always wins! (P = Picard)

## III Quantum Prisoner's Dilemma

Rule: No communication btwn players

Def:  $|c\rangle, |d\rangle$  Hilbert space basis



$$J = \frac{1}{\sqrt{2}} (JJ + iXX)$$

$$= \exp(i\pi/4 XX)$$

Choose  $\hat{C} := I$   
 $\hat{D} := X$  } choices for  $U_A, U_B$

$\Rightarrow$  Game reduces to classical game =

$$0 = [J, \hat{C}\hat{C}] = [J, \hat{C}\hat{D}] = [J, \hat{D}\hat{C}] = [J, \hat{D}\hat{D}]$$

Measure prob (CC, CD, DC, DD)

Payoff:

$$\begin{cases} \$A = r P_{CC} + p P_{DD} + t P_{DC} + s' P_{CD} \\ \$B = r P_{CC} + p P_{DD} + t P_{CD} + s' P_{DC} \end{cases}$$

reward
punishment
temptation
suicider's payoff

let  $r=3, p=1, t=5, s=0$  as before

New quantum strategy:  $\hat{Q} = Z$

		$u_B$		
		C	D	Q
$u_A$	C	3,3	0,5	-1,1
	D	5,0	1,1	0,5
	Q	1,1	5,0	3,3

Shaded:  
New!  
(See below)

$u_A u_B$

$$\begin{aligned} XZX &= -Z \\ ZX &= iY \end{aligned}$$

$$\begin{aligned} \hat{Q} \hat{C} : |\Psi\rangle &= J^\dagger (Z \otimes I) J |CC\rangle \\ &= (II - iXX)(ZI)(II + iXX) |CC\rangle \\ &= (ZI - YX - YX - ZI) |CC\rangle \\ &= -2YX |CC\rangle \\ &\approx |DD\rangle \text{ (when normalized)} \end{aligned}$$

$$\begin{aligned} \hat{Q} \hat{Q} : |\Psi\rangle &= (II - iXX)(ZZ)(II + iXX) |CC\rangle \\ &= |CC\rangle \end{aligned}$$

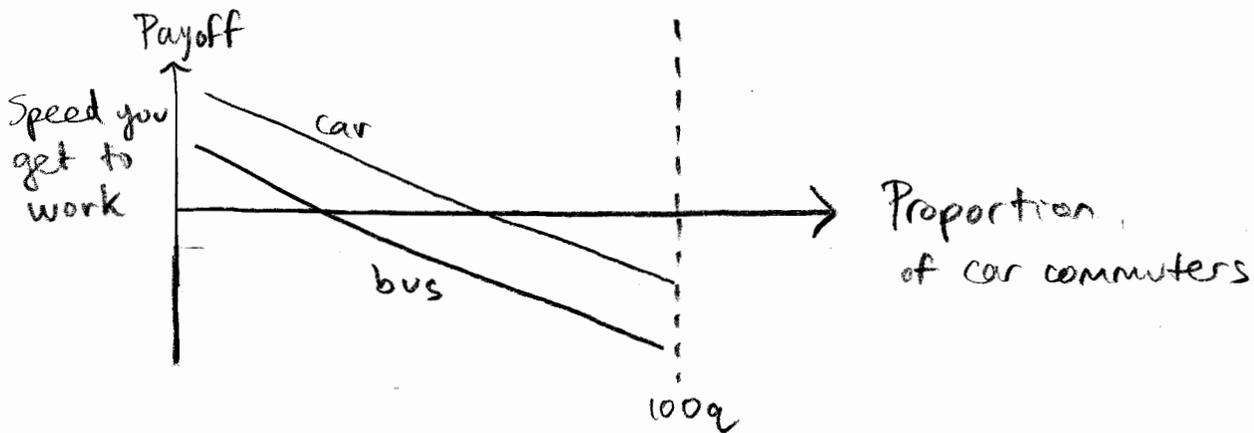
Old Nash equil: DD

New Nash equil: QQ (better!)

- Problems:
- Why help prisoners by offering  $\varphi$  choice
  - Who enforces  $I$ ?
  - What about other rotations?

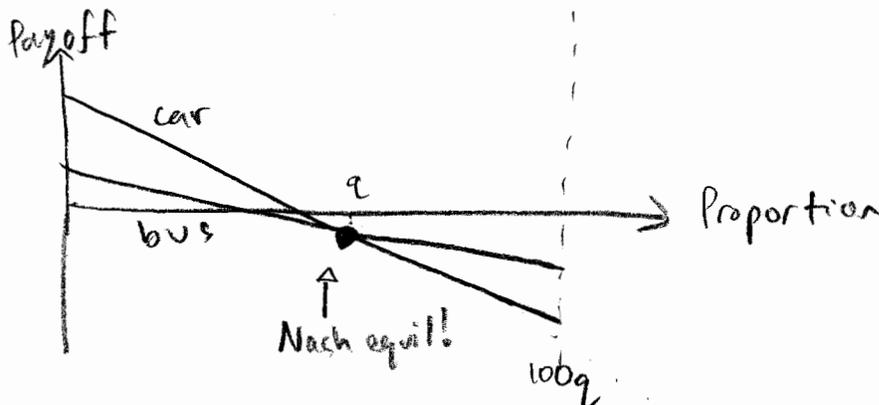
## IV Tragedy of the Commons

Proportion game model:



Dominant strategy: Take car. Sucks!

New model:



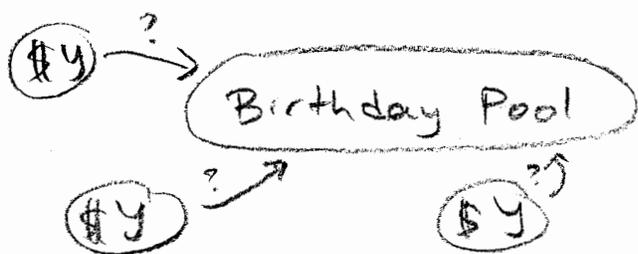
e.g. Govt provides incentives to take bus till get interests

Tragedy: All common property resources tend to be overexploited & thus degraded.

Traditional sol<sup>n</sup>: 3rd-party regulator (e.g. govt)

IV

Q Public Goods Game



$n$  = # players  
 $y$  = initial endowment  
 $C_k$  = contribution from player  $k$   
 $a$  = public gain,  $1 < a < n$

$$\$k = y - C_k + \frac{a}{n} \sum_j C_j$$

Claim: Nash equilib is  $\sum_j C_j = 0$

let  $y=1, n=2$ :

		Alice	
		C	D
Alice	C	$a, a$	$a/2, 1+a/2$
	D	$1+a/2, a$	$1, 1$

Same matrix as prisoner's dilemma! Oh no!

No one gives to birthday pool, though it'd be awesome if everyone did

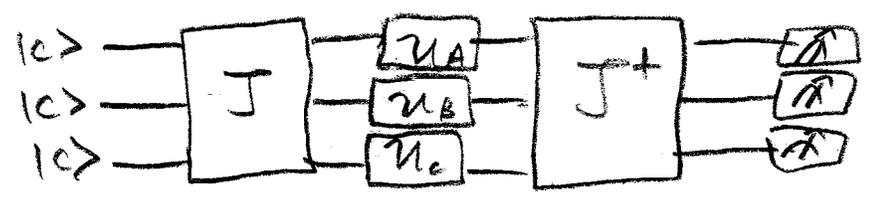
let  $n=3$ , 0=defect, 1=contribute

	ABC	\$A	\$B	\$C	$a \sum C_i$
No one contributes	000	1	1	1	0
Only 1 contributes	001	$1+a/3$	$1+a/3$	$a/3$	$a$
	010	$1+a/3$	$1+a/3$	$a/3$	$a$
2 contribute	011	$1+2a/3$	$2a/3$	$2a/3$	$2a$
	100	$1+2a/3$	$2a/3$	$2a/3$	$2a$
Everyone contributes	101	$a$	$a$	$a$	$3a$
	110	$a$	$a$	$a$	$3a$
	111	$a$	$a$	$a$	$3a$

Nash equil: 000 (sad!)

Quantum Game

chen, Hogg, ...  
quant-ph/0301013



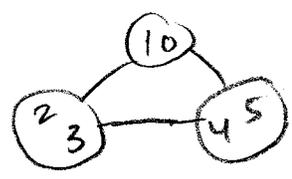
Assume  $J$  &  $J^+$  performed by trusted authority (eg. bank)

- let  $J = \frac{1}{\sqrt{2}} (III + iXXX)$

Claim: Expected payoff =  $\frac{1}{2} (1+a)$

- let  $J$  have pairwise entanglement

Claim: Expected payoff =  $a - 2^{-(n-1)} (n-1) \approx a$



Problems: Will people trust J? Trust quantum mech.?

Will people collude?

Will people swap in their own qstates?