Selinger Optimizer

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(Slides gently offered by Sam Madden)

The Problem

How to order a series of N joins, e.g.,

$$A.a = B.b AND A.c = D.d AND B.e = C.f$$

N! ways to order joins (e.g., ABCD, ACBD,)
(N-1)! plans per ordering (e.g., (((AB)C)D), ((AB)(CD), ...)
Multiple implementations (e.g., hash, nested loops, etc)

- Naïve approach doesn't scale, e.g., for 20-way join
 - $-10! \times 9! = 1.3 \times 10^{12}$
 - $-20! \times 19! = 2.9 \times 10^{35}$

Selinger Optimizations

- Left-deep only (((AB)C)D) (eliminate (N-1)!)
- Push-down selections
- Don't consider cross products
- Dynamic programming algorithm

Dynamic Programming

```
R \leftarrow set of relations to join (e.g., ABCD)

For \partial in \{1...|R|\}:

for S in \{\text{all length } \partial subsets of R\}:

optjoin(S) = a join (S-a),

where a is the single relation that minimizes:

\text{cost}(\text{optjoin}(S-a)) +

min. cost to join (S-a) to a +

min. access cost for a
```

optjoin(S-a) is cached from previous iteration

Cache Subplan Best Cost choice A index 100 B seq scan 50 ...

optjoin(ABCD) – assume all joins are NL

0=1

A = best way to access A

(e.g., sequential scan, or predicate pushdown into index...)

B = best way to access B

C = best way to access C

D = best way to access D

Total cost computations: *choose*(N,1), where N is number of relations

Cache

Cache				
Subplan	Best choice	Cost		
А	index	100		
В	seq scan	50		
{A,B}	ВА	156		
{B,C}	ВС	98		
•••				

optjoin(ABCD)

$$\{A,B\} = AB \text{ or } BA$$

(using previously computed best way to access A and B)

$$\{B,C\} = BC \text{ or } CB$$

$$\{C,D\} = CD \text{ or } DC$$

$$\{A,C\} = AC \text{ or } CA$$

$$\{A,D\} = AD \text{ or } DA$$

$$\{B,D\} = BD \text{ or } DB$$

Total cost computations: choose(N,2) x 2

optjoin(ABCD)

${A,C,D} = ...$

$$\{B,C,D\} = ...$$

Cache

Subplan	Best choice	Cost
Α	index	100
В	seq scan	50
{A,B}	ВА	156
{B,C}	ВС	98
{A,B,C}	ВСА	125
{B,C,D}	BCD	115

Total cost computations: $choose(N,3) \times 3 \times 2$

Already computed − lookup in cache {A,B,C,D} = remove A, compare A {B,C,D}) to ({B,C,D})A remove B, compare B({A,C,D}) to ({A,C,D})B remove C, compare C({A,B,D}) to ({A,B,D})C remove D, compare D({A,B,C}) to ({A,B,C})D

Final answer is plan with minimum cost of these four

Total cost computations: $choose(N,4) \times 4 \times 2$

Cache

Subplan	Best choice	Cost
Α	index	100
В	seq scan	50
{A,B}	ВА	156
{B,C}	ВС	98
{A,B,C}	ВСА	125
{B,C,D}	BCD	115
{A,B,C,D}	ABCD	215

Complexity

choose(n,1) + choose(n,2) + ... + choose(n,n) total
subsets considered

All subsets of a size n set = power set of $n = 2^n$

Equiv. to computing all binary strings of size n 000,001,010,100,011,101,110,111

Each bit represents whether an item is in or out of set

Complexity (continued)

For each subset,

k ways to remove 1 join

k < n

m ways to join 1 relation with remainder

Total cost: O(nm2^n) plan evaluations n = 20, m = 2 4.1×10^7

Interesting Orders

- Some queries need data in sorted order
 - Some plans produce sorted data (e.g., using an index scan or merge join
- May be non-optimal way to join data, but overall optimal plan
 - Avoids final sort
- In cache, maintain best overall plan, plus best plan for each interesting order
- At end, compare cost of

```
best plan + sort into ordertobest in order plan
```

 Increases complexity by factor of k+1, where k is number of interesting orders

SELECT A.f3, B.f2 FROM A,B where A.f3 = B.f4 ORDER BY A.f3

Subplan	Best choice	Cost	Best in A.f3 order	Cost
Α	index	100	index	100
В	seq scan	50	seqscan	50
{A,B}	BA hash	156	AB merge	180

```
compare:
    cost(sort(output)) + 156
to
180
```

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