

6.241: Dynamic Systems—Fall 2007

HOMEWORK 3 SOLUTIONS

Exercise 3.2 i) We would like to minimize the 2-norm of u , i.e., $\|\underline{u}\|_2^2$. Since y_n is given as

$$y_n = \sum_{i=1}^n h_i u_{n-1}$$

we can rewrite this equality as

$$y_n = \begin{bmatrix} h_1 & h_2 & \cdots & h_n \end{bmatrix} \begin{bmatrix} u_{n-1} \\ u_{n-2} \\ \vdots \\ u_0 \end{bmatrix}$$

We want to find the \underline{u} with the smallest 2-norm such that

$$\bar{y} = A\underline{u}.$$

where we assume that A has a full rank (i.e. $h_i \neq 0$ for some i , $1 \leq i \leq n$). Then, the solution reduces to the familiar form:

$$\hat{u} = A'(AA')^{-1}\bar{y}.$$

By noting that $AA' = \sum_{i=1}^n h_i^2$, we can obtain \hat{u}_j as follows;

$$\hat{u}_j = \frac{h_j \bar{y}}{\sum_{i=1}^n h_i^2}, \quad \text{for } j = 0, 1, \dots, n-1.$$

ii) a) Let's introduce e as an error such that $y_n = \bar{y} - e$. It can also be written as $\bar{y} - y_n = e$. Then now the quantity we would like to minimize can be written as

$$r(\bar{y} - y_n)^2 + u_0^2 + \cdots + u_{n-1}^2$$

where r is a positive weighting parameter. The problem becomes to solve the following minimization problem :

$$\hat{u} = \arg \min_u \sum_{i=1}^n u_i^2 + re^2 = \arg \min_u (\|\underline{u}\|_2^2 + r\|e\|_2^2),$$

from which we see that r is a weight that characterizes the tradeoff between the size of the final error, $\bar{y} - y_n$, and energy of the input signal, \underline{u} .

In order to reduce the problem into the familiar form, i.e., $\|y - Ax\|$, let's augment $\sqrt{r}e$ at the bottom of \underline{u} so that a new augmented vector, \tilde{u} is

$$\tilde{u} = \begin{bmatrix} \underline{u} \\ \sqrt{r}e \end{bmatrix},$$

This choice of \tilde{u} follows from the observation that this is the \tilde{u} that would have $\|\tilde{u}\|_2^2 = \|\underline{u}\|_2^2 + re^2$, the quantity we aim to minimize .

Now we can write \bar{y} as follows

$$\bar{y} = \begin{bmatrix} A & \vdots & \frac{1}{\sqrt{r}} \end{bmatrix} \begin{bmatrix} \underline{u} \\ \cdots \\ \sqrt{r}e \end{bmatrix} = \tilde{A}\tilde{u} = A\underline{u} + e = y_n + e.$$

Now, \hat{u} can be obtained using the augmented A , \tilde{A} , as

$$\hat{u} = \tilde{A}'(\tilde{A}\tilde{A}')^{-1}\bar{y} = \begin{bmatrix} A' \\ \frac{1}{\sqrt{r}} \end{bmatrix} \begin{bmatrix} AA' + \frac{1}{r} \end{bmatrix} \bar{y}.$$

By noting that

$$AA' + \frac{1}{r} = \sum_{i=1}^n h_i^2 + \frac{1}{r},$$

we can obtain \hat{u}_j as follows

$$\hat{u}_j = \frac{h_j \bar{y}}{\sum_{i=1}^n h_i^2 + \frac{1}{r}} \quad \text{for } j = 0, \dots, n-1.$$

ii) b) When $r = 0$, it can be interpreted that the error can be anything, but we would like to minimize the input energy. Thus we expect that the solution will have all the u_i 's to be zero. In fact, the expression obtained in ii) a) will be zero as $r \rightarrow 0$. On the other hand, the other situation is an interesting case. We put a weight of ∞ to the final state error, then the expression from ii) a) gives the same expression as in i) as $r \rightarrow \infty$.

Exercise 3.3 This problem is similar to Example 3.4, except now we require that $\dot{p}(T) = 0$. We can derive, from $x(t) = \ddot{p}(t)$, that $p(t) = x(t) * tu(t) = \int_0^t (t-\tau)x(\tau)d\tau$ where $*$ denotes convolution and $u(t)$ is the unit step, defined as 1 when $t > 0$ and 0 when $t < 0$. (One way to derive this is to take $x(t) = \ddot{p}(t)$ to the Laplace domain, *taking into account initial conditions*, to find the transfer function $H(s) = P(s)/X(s)$, hence the impulse response, $h(t)$ such that $p(t) = x(t) * h(t)$). Similarly, $\dot{p}(t) = x(t) * u(t) = \int_0^t x(\tau)d\tau$. So, $y = p(T) = \int_0^T (T-\tau)x(\tau)d\tau$ and $0 = \dot{p}(T) = x(T) * u(T) = \int_0^T x(\tau)d\tau$. You can check that $\langle g(t), f(t) \rangle = \int_0^T g(t)f(t)d\tau$ is an inner product on the space of continuous functions on $[0, T]$, denoted by $C[0, T]$, which we are searching for $x(t)$. So, we have that $y = p(T) = \langle (T-t), x(t) \rangle$ and $0 = \dot{p}(T) = \langle 1, x(t) \rangle$. In matrix form,

$$\begin{bmatrix} y \\ 0 \end{bmatrix} = \begin{bmatrix} \langle T-t, x(t) \rangle \\ \langle 1, x(t) \rangle \end{bmatrix} = \prec [T-t \quad 1], x(t) \succ$$

where $\prec \dots \succ$ denotes the Grammian, as defined in chapter 2. Now, in chapter 3, it was shown that the minimum length solution to $y = \prec A, x \succ$, is $\hat{x} = A \prec A, A \succ^{-1} y$. So, for our problem,

$$\hat{x} = [T-t \quad 1] \prec [T-t \quad 1], [T-t \quad 1] \succ^{-1} \begin{bmatrix} y \\ 0 \end{bmatrix}.$$

Where, using the definition of the Grammian, we have that:

$$\prec [T-t \quad 1], [T-t \quad 1] \succ = \begin{bmatrix} \langle T-t, T-t \rangle & \langle T-t, 1 \rangle \\ \langle 1, T-t \rangle & \langle 1, 1 \rangle \end{bmatrix}.$$

Now, we can use the definition for inner product to find the individual entries, $\langle T-t, T-t \rangle = \int_0^T (T-t)^2 dt = T^3/3$, $\langle T-t, 1 \rangle = \int_0^T (T-t) dt = T^2/2$, and $\langle 1, 1 \rangle = T$. Plugging these in, one can simplify the expression for \hat{x} and obtain $\hat{x}(t) = \frac{12y}{T^2} [\frac{1}{2} - \frac{t}{T}]$ for $t \in [0, T]$.

Alternatively, we have that $x(t) = \ddot{p}(t)$. Integrating both sides and taking into account that $p(0) = 0$ and $\dot{p}(0) = 0$, we have $p(t) = \int_0^t \int_0^{t_1} x(\tau) d\tau dt_1 = \int_0^t f(t_1) dt_1$. Now, we use the integration by parts formula, $\int_0^t u dv = uv|_0^t - \int_0^t v du$, with $u = f(t_1) = \int_0^{t_1} x(\tau) d\tau$, and $dv = dt_1$; hence $du = df(t_1) = x(t_1) dt_1$ and $v = t_1$. Plugging in and simplifying we get that $p(t) = \int_0^t \int_0^{t_1} x(\tau) d\tau dt_1 = \int_0^t (t-\tau)x(\tau) d\tau$. Thus, $y = p(T) = \int_0^T (T-\tau)x(\tau) d\tau = \langle T-t, x(t) \rangle$. In addition, we have that $0 = \dot{p}(T) = \int_0^T x(\tau) d\tau = \langle 1, x(t) \rangle$. That is, we seek to find the minimum length $x(t)$ such that

$$\begin{aligned} y &= \langle T-t, x(t) \rangle \\ 0 &= \langle 1, x(t) \rangle. \end{aligned}$$

Recall that the minimum length solution $\hat{x}(t)$ must be a linear combination of $T-t$ and 1 , i.e., $\hat{x}(t) = a_1(T-t) + a_2$. So,

$$\begin{aligned} y &= \langle T-t, a_1(T-t) + a_2 \rangle = a_1 \int_0^T (T-t)^2 dt + a_2 \int_0^T (T-t) dt = a_1 \frac{T^3}{3} + a_2 \frac{T^2}{2} \\ 0 &= \langle 1, a_1(T-t) + a_2 \rangle = \int_0^T (a_1(T-t) + a_2) dt = a_1 \frac{T^2}{2} + a_2 T. \end{aligned}$$

This is a system of two equations and two unknowns, which we can rewrite in matrix form:

$$\begin{bmatrix} y \\ 0 \end{bmatrix} = \begin{bmatrix} \frac{T^3}{3} & \frac{T^2}{2} \\ \frac{T^2}{2} & T \end{bmatrix} \begin{bmatrix} a_1 \\ a_2 \end{bmatrix},$$

So,

$$\begin{bmatrix} a_1 \\ a_2 \end{bmatrix} = \begin{bmatrix} \frac{T^3}{3} & \frac{T^2}{2} \\ \frac{T^2}{2} & T \end{bmatrix}^{-1} \begin{bmatrix} y \\ 0 \end{bmatrix}.$$

Exercise 4.1 Note that for any $v \in C^m$, (show this!)

$$\|v\|_\infty \leq \|v\|_2 \leq \sqrt{m} \|v\|_\infty. \quad (1)$$

Therefore, for $A \in C^{m \times n}$ with $x \in \mathbb{C}^n$

$$\|Ax\|_2 \leq \sqrt{m} \|Ax\|_\infty \rightarrow \text{for } x \neq 0, \frac{\|Ax\|_2}{\|x\|_2} \leq \sqrt{m} \frac{\|Ax\|_\infty}{\|x\|_2}.$$

But, from equation (1), we also know that $\frac{1}{\|x\|_\infty} \geq \frac{1}{\|x\|_2}$. Thus,

$$\frac{\|Ax\|_2}{\|x\|_2} \leq \frac{\sqrt{m} \|Ax\|_\infty}{\|x\|_2} \leq \frac{\sqrt{m} \|Ax\|_\infty}{\|x\|_\infty} \leq \sqrt{m} \|A\|_\infty, \quad (2)$$

Equation (2) must hold for all $x \neq 0$, therefore

$$\max_{x \neq 0} \frac{\|Ax\|_2}{\|x\|_2} = \|A\|_2 \leq \sqrt{m}\|A\|_\infty.$$

To prove the lower bound $\frac{1}{\sqrt{n}}\|A\|_\infty \leq \|A\|_2$, reconsider equation (1):

$$\|Ax\|_\infty \leq \|Ax\|_2 \rightarrow \text{for } x \neq 0, \frac{\|Ax\|_\infty}{\|x\|_2} \leq \frac{\|Ax\|_2}{\|x\|_2} \leq \|A\|_2 \rightarrow \frac{\sqrt{n}\|Ax\|_\infty}{\|x\|_2} \leq \frac{\sqrt{n}\|Ax\|_2}{\|x\|_2} \leq \sqrt{n}\|A\|_2. \quad (3)$$

But, from equation (1) for $x \in \mathbb{C}^n$, $\frac{\sqrt{n}}{\|x\|_2} \geq \frac{1}{\|x\|_\infty}$. So,

$$\frac{\|Ax\|_\infty}{\|x\|_\infty} \leq \frac{\sqrt{n}\|Ax\|_\infty}{\|x\|_2} \leq \sqrt{n}\|A\|_2$$

for all $x \neq 0$ including x that makes $\frac{\|Ax\|_\infty}{\|x\|_\infty}$ maximum, so,

$$\max_{x \neq 0} \frac{\|Ax\|_\infty}{\|x\|_\infty} = \|A\|_\infty \leq \sqrt{n}\|A\|_2,$$

or equivalently,

$$\frac{1}{\sqrt{n}}\|A\|_\infty \leq \|A\|_2.$$

Exercise 4.5 Any $m \times n$ matrix A , it can be expressed as

$$A = U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V',$$

where U and V are unitary matrices. The ‘‘Moore-Penrose inverse’’, or *pseudo-inverse* of A , denoted by A^+ , is then defined as the $n \times m$ matrix

$$A^+ = V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U'.$$

a) Now we have to show that A^+A and AA^+ are symmetric, and that $AA^+A = A$ and $A^+AA^+ = A^+$. Suppose that Σ is a diagonal invertible matrix with the dimension of $r \times r$. Using the given definitions as well as the fact that for a unitary matrix U , $U'U = UU' = I$, we have

$$\begin{aligned} AA^+ &= U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V'V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U' \\ &= U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} I \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U' \\ &= U \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} U', \end{aligned}$$

which is symmetric. Similarly,

$$\begin{aligned}
A^+A &= V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U'U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V' \\
&= V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} I \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V' \\
&= V \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} V'
\end{aligned}$$

which is again symmetric.

The facts derived above can be used to show the other two.

$$\begin{aligned}
AA^+A &= (AA^+)A = U \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} U'A \\
&= U \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} U'U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V' \\
&= U \begin{pmatrix} \Sigma & 0 \\ 0 & 0 \end{pmatrix} V' \\
&= A.
\end{aligned}$$

Also,

$$\begin{aligned}
A^+AA^+ &= (A^+A)A^+ = V \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} V'A^+ \\
&= V \begin{pmatrix} I_{r \times r} & 0 \\ 0 & 0 \end{pmatrix} V'V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U' \\
&= V \begin{pmatrix} \Sigma^{-1} & 0 \\ 0 & 0 \end{pmatrix} U' \\
&= A^+.
\end{aligned}$$

b) We have to show that when A has full column rank then $A^+ = (A'A)^{-1}A'$, and that when A has full row rank then $A^+ = A'(AA')^{-1}$. If A has full column rank, then we know that $m \geq n$, $\text{rank}(A) = n$, and

$$A = U \begin{pmatrix} \Sigma_{n \times n} \\ 0 \end{pmatrix} V'.$$

Also, as shown in chapter 2, when A has full column rank, $(A'A)^{-1}$ exists. Hence

$$\begin{aligned}
(A'A)^{-1}A' &= \left(V \begin{pmatrix} \Sigma' & 0 \end{pmatrix} U'U \begin{pmatrix} \Sigma \\ 0 \end{pmatrix} V' \right)^{-1} V \begin{pmatrix} \Sigma' & 0 \end{pmatrix} U' \\
&= (V\Sigma'\Sigma V')^{-1} V \begin{pmatrix} \Sigma' & 0 \end{pmatrix} U' \\
&= V(\Sigma'\Sigma)^{-1} V'V \begin{pmatrix} \Sigma' & 0 \end{pmatrix} U' \\
&= V(\Sigma'\Sigma)^{-1} \begin{pmatrix} \Sigma' & 0 \end{pmatrix} U' \\
&= V \begin{pmatrix} \Sigma^{-1} & 0 \end{pmatrix} U' \\
&= A^+.
\end{aligned}$$

Similarly, if A has full row rank, then $n \geq m$, $\text{rank}(A) = m$, and

$$A = U \begin{pmatrix} \Sigma_{m \times m} & 0 \end{pmatrix} V'.$$

It can be proved that when A has full row rank, $(A'A)^{-1}$ exists. Hence,

$$\begin{aligned} A'(AA')^{-1} &= V \begin{pmatrix} \Sigma' \\ 0 \end{pmatrix} U' \left(U \begin{pmatrix} \Sigma & 0 \end{pmatrix} V' V \begin{pmatrix} \Sigma' \\ 0 \end{pmatrix} U' \right)^{-1} \\ &= V \begin{pmatrix} \Sigma' \\ 0 \end{pmatrix} U' (U \Sigma \Sigma' U')^{-1} \\ &= V \begin{pmatrix} \Sigma' \\ 0 \end{pmatrix} U' U (\Sigma \Sigma^{-1}) U' \\ &= V \begin{pmatrix} \Sigma^{-1} \\ 0 \end{pmatrix} U' \\ &= A^+. \end{aligned}$$

c) Show that, of all x that minimize $\|y - Ax\|_2$, the one with the smallest length $\|x\|_2$ is given by $\hat{x} = A^+y$. If A has full row rank, we have shown in chapter 3 that the solution with the smallest length is given by

$$\hat{x} = A'(AA')^{-1}y,$$

and from part (b), $A'(AA')^{-1} = A^+$. Therefore

$$\hat{x} = A^+y.$$

Similarly, it can be shown that the pseudo inverse is the solution for the case when a matrix A has a full column rank (compare the results in chapter 2 with the expression you found in part (b) for A^+ when A has full column rank).

Now, let's consider the case when a matrix A is rank deficient, *i.e.*, $\text{rank}(A) = r < \min(m, n)$ where $A \in C^{m \times n}$ and is thus neither full row or column rank. Suppose we have a singular value decomposition of A as

$$A = U \Sigma V',$$

where U and V are unitary matrices. Then the norm we are minimizing is

$$\|Ax - y\| = \|U \Sigma V' x - y\| = \|U(\Sigma V' x - U' y)\| = \|\Sigma z - U' y\|,$$

where $z = V'x$, since $\|\cdot\|$ is unaltered by the orthogonal transformation, U . Thus, x minimizes $\|Ax - y\|$ if and only if z minimizes $\|\Sigma z - c\|$, where $c = U' y$. Since the rank of A is r , the matrix Σ has the nonzero singular values $\sigma_1, \sigma_2, \dots, \sigma_r$ in its diagonal entries. Then we can rewrite $\|\Sigma z - c\|^2$ as follows:

$$\|\Sigma z - c\|^2 = \sum_{i=1}^r (\sigma_i z_i - c_i)^2 + \sum_{i=r+1}^n c_i^2.$$

It is clear that the minimum of the norm can be achieved when $z_i = \frac{c_i}{\sigma_i}$ for $i = 1, 2, \dots, r$ and the rest of the z_i 's can be chosen arbitrarily. Thus, there are infinitely many solutions \hat{z} and the

solution with the minimum norm can be achieved when $z_i = 0$ for $i = r + 1, r + 2, \dots, n$. Thus, we can write this \hat{z} as

$$z = \Sigma_1 c,$$

where

$$\Sigma_1 = \begin{pmatrix} \Sigma_r^{-1} & 0 \\ 0 & 0 \end{pmatrix}$$

and Σ_r is a square matrix with nonzero singular values in its diagonal in decreasing order. This value of z also yields the value of x of minimal 2 norm since V is a unitary matrix.

Thus the solution to this problem is

$$\hat{x} = Vz = V\Sigma_1 c = V\Sigma_1 U' y = A^+ y.$$

It can be easily shown that this choice of A^+ satisfies all the conditions, or definitions, of pseudo inverse in a).

Exercise 4.6. a) Suppose $A \in C_n^m$ has full column rank. Then QR factorization for A can be easily constructed from SVD:

$$A = U \begin{pmatrix} \Sigma_n \\ 0 \end{pmatrix} V'$$

where Σ_n is a $n \times n$ diagonal matrix with singular values on the diagonal. Let $Q = U$ and $R = \Sigma_n V'$ and we get the QR factorization. Since Q is an orthogonal matrix, we can represent any $Y \in C_p^m$ as

$$Y = Q \begin{pmatrix} Y_1 \\ Y_2 \end{pmatrix}$$

Next

$$\|Y - AX\|_F^2 = \|Q \begin{pmatrix} Y_1 \\ Y_2 \end{pmatrix} - Q \begin{pmatrix} R \\ 0 \end{pmatrix} X\|_F^2 = \|Q \begin{pmatrix} Y_1 - RX \\ Y_2 \end{pmatrix}\|_F^2$$

Denote

$$D = \begin{pmatrix} Y_1 - RX \\ Y_2 \end{pmatrix}$$

and note that multiplication by an orthogonal matrix does not change Frobenius norm of the matrix:

$$\|QD\|_F^2 = \text{tr}(D'Q'QD) = \text{tr}(D'D) = \|D\|_F^2$$

Since Frobenius norm squared is equal to sum of squares of all elements, square of the Frobenius norm of a block matrix is equal to sum of the squares of Frobenius norms of the blocks:

$$\left\| \begin{pmatrix} Y_1 - RX \\ Y_2 \end{pmatrix} \right\|_F^2 = \|Y_1 - RX\|_F^2 + \|Y_2\|_F^2$$

Since Y_2 block can not be affected by choice of X matrix, the problem reduces to minimization of $\|Y_1 - RX\|_F^2$. Recalling that R is invertible (because A has full column rank) the solution is

$$X = R^{-1}Y_1$$

b) Evaluate the expression with the pseudoinverse using the representations of A and Y from part a):

$$(A'A)^{-1} A'Y = \left([R' \ 0] Q'Q \begin{bmatrix} R \\ 0 \end{bmatrix} \right)^{-1} [R' \ 0] Q'Q \begin{bmatrix} Y_1 \\ Y_2 \end{bmatrix} = R^{-1} (R')^{-1} [R' \ 0] \begin{bmatrix} Y_1 \\ Y_2 \end{bmatrix} = R^{-1} Y_1$$

From 4.5 b) we know that if a matrix has a full column rank, $A^+ = (A'A)^{-1} A'$, therefore both expressions give the same solutions.

c)

$$\|Y - AX\|_F^2 + \|Z - BX\|_F^2 = \left\| \begin{pmatrix} Y \\ Z \end{pmatrix} - \begin{pmatrix} A \\ B \end{pmatrix} X \right\|_F^2$$

Since A has full column rank, $\begin{pmatrix} A \\ B \end{pmatrix}$ also has full column rank, therefore we can apply results from parts a) and b) to conclude that

$$X = \left(\begin{pmatrix} A \\ B \end{pmatrix}' \begin{pmatrix} A \\ B \end{pmatrix} \right)^{-1} \begin{pmatrix} A \\ B \end{pmatrix}' \begin{pmatrix} Y \\ Z \end{pmatrix} = (A'A + B'B)^{-1} (A'Y + B'Z)$$

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Spring 2011

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