

[Show All Answers](#)[Hide All Answers](#)

6.02 Practice Problems: Error Correcting Codes

Problem 1. For each of the following sets of codewords, please give the appropriate (n,k,d) designation where n is number of bits in each codeword, k is the number of message bits transmitted by each code word and d is the minimum Hamming distance between codewords. Also give the code rate.

A. {111, 100, 001, 010}

[Show Answer](#)

B. {00000, 01111, 10100, 11011}

[Show Answer](#)

C. {00000}

[Show Answer](#)

Problem 2. Suppose management has decided to use 20-bit data blocks in the company's new (n,20,3) error correcting code. What's the minimum value of n that will permit the code to be used for single bit error correction?

[Show Answer](#)

Problem 3. The Registrar has asked for an encoding of class year ("Freshman", "Sophomore", "Junior", "Senior") that will allow single error correction. Please give an appropriate 5-bit binary encoding for each of the four years.

[Show Answer](#)

Problem 4. For any block code with minimum Hamming distance at least $2t + 1$ between code words, show that:

$$2^{n-k} \geq 1 + \binom{n}{1} + \binom{n}{2} + \dots + \binom{n}{t}.$$

[Show Answer](#)

Problem 5. Pairwise Communications has developed a block code with three data (D1, D2, D3) and three parity bits (P1, P2, P3):

$$P1 = D1 + D2$$

$$P2 = D2 + D3$$

$$P3 = D3 + D1$$

A. What is the (n,k,d) designation for this code.

Show Answer

B. The receiver computes three syndrome bits from the (possibly corrupted) received data and parity bits:

$$\begin{aligned}
E1 &= D1 + D2 + P1 \\
E2 &= D2 + D3 + P2 \\
E3 &= D3 + D1 + P3.
\end{aligned}$$

The receiver performs maximum likelihood decoding using the syndrome bits. For the combinations of syndrome bits listed below, state what the maximum-likelihood decoder believes has occurred: no errors, a single error in a specific bit (state which one), or multiple errors.

$$\begin{aligned}
E1 \ E2 \ E3 &= 000 \\
E1 \ E2 \ E3 &= 010 \\
E1 \ E2 \ E3 &= 101 \\
E1 \ E2 \ E3 &= 111
\end{aligned}$$

Show Answer

Problem 6. Dos Equis Encodings, Inc. specializes in codes that use 20-bit transmit blocks. They are trying to design a (20, 16) linear block code for single error correction. Explain whether they are likely to succeed or not.

Show Answer

Problem 7. Consider the following (n,k,d) block code:

D0	D1	D2	D3	D4		P0
D5	D6	D7	D8	D9		P1
D10	D11	D12	D13	D14		P2
P3	P4	P5	P6	P7		

where D0-D14 are data bits, P0-P2 are row parity bits and P3-P7 are column parity bits. The transmitted code word will be:

D0 D1 D2 ... D13 D14 P0 P1 ... P6 P7

A. Please give the values for n, k, d for the code above.

Show Answer

B. If $D_0 \ D_1 \ D_2 \ \dots \ D_{13} \ D_{14} = 0 \ 1 \ 0 \ 1 \ 0, \ 0 \ 1 \ 0 \ 0 \ 1, \ 1 \ 0 \ 0 \ 0 \ 1$, please compute P0 through P7.

Show Answer

C. Now we receive the four following code words:

M1: 0 1 0 1 0, 0 1 0 0 1, 1 0 0 0 1, 0 0 0 1 1 0 1 0
M2: 0 1 0 1 0, 0 1 0 0 1, 1 0 0 0 1, 0 0 1 1 1 0 1 0
M3: 0 1 0 1 0, 0 1 0 0 1, 1 0 0 0 1, 1 1 0 1 1 0 1 0
M4: 0 1 0 1 0, 0 1 0 0 1, 1 0 0 0 1, 1 0 0 1 1 0 1 0

For each of received code words, indicate the number of errors. If there are errors, indicate if they are

correctable, and if they are, what the correction should be.

Show Answer

Problem 8. The following matrix shows a rectangular single error correcting code consisting of 9 data bits, 3 row parity bits and 3 column parity bits. For each of the examples that follow, please indicate the correction the receiver must perform: give the position of the bit that needs correcting (e.g., D7, R1), or "no" if there are no errors, or "M" if there is a multi-bit uncorrectable error.

D1	D2	D3	R1
D4	D5	D6	R2
D7	D8	D9	R3
C1	C2	C3	—

1	1	1	1	0	0	1	0	1	1	0	0	1	0	1	0	1	1	1	0
1	1	0	0	0	1	1	1	1	0	0	1	1	0	0	1	1	1	1	1
0	1	1	0	0	1	1	0	0	0	1	0	1	1	0	0	1	0	0	0
0	1	1	—	0	0	1	—	0	0	1	—	1	1	1	—	0	0	1	—

Show Answer

Problem 9. Consider two convolutional coding schemes - I and II. The generator polynomials for the two schemes are

Scheme I: $G_0 = 1101$, $G_1 = 1110$

Scheme II: $G_0 = 110101$, $G_1 = 111011$

Notation is follows: if the generator polynomial is, say, 1101, then the corresponding parity bit for message bit n is

$$(x[n] + x[n-1] + x[n-3]) \bmod 2$$

where $x[n]$ is the message sequence.

- A. Indicate TRUE or FALSE
 - a. Code rate of Scheme I is 1/4.
 - b. Constraint length of Scheme II is 4.
 - c. Code rate of Scheme II is equal to code rate of Scheme I.
 - d. Constraint length of Scheme I is 4.

Show Answer

- B. How many states will there be in the state diagram for Scheme I? For Scheme II?

Show Answer

- C. Which code will lead to a lower bit error rate? Why?

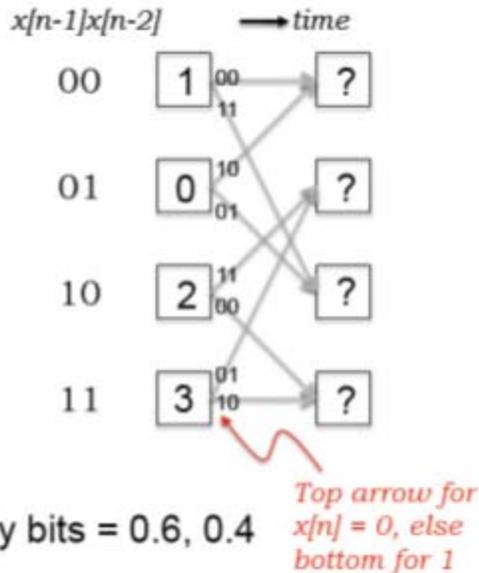
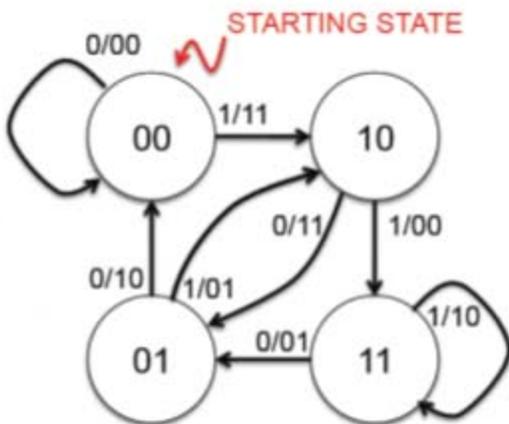
Show Answer

- D. Alyssa P. Hacker suggests a modification to Scheme I which involves adding a third generator

polynomial $G_2 = 1001$. What is the code rate r of Alyssa's coding scheme? What about constraint length κ ? Alyssa claims that her scheme is stronger than Scheme I. Based on your computations for r and κ , is her statement true?

Show Answer

Problem 10. Consider a convolution code that uses two generator polynomials: $G_0 = 111$ and $G_1 = 110$. You are given a particular snapshot of the decoding trellis used to determine the most likely sequence of states visited by the transmitter while transmitting a particular message:



A. Complete the Viterbi step, i.e., fill in the question marks in the matrix, assuming a hard branch metric based on the Hamming distance between expected and received parity where the received voltages are digitized using a 0.5V threshold.

Show Answer

B. Complete the Viterbi step, i.e., fill in the question marks in the matrix, assuming a soft branch metric based on the square of the Euclidean distance between expected and received parity voltages. Note that your branch and path metrics will not necessarily be integers.

Show Answer

C. Does the soft metric give a different answer than the hard metric? Base your response in terms of the relative ordering of the states in the second column and the survivor paths.

Show Answer

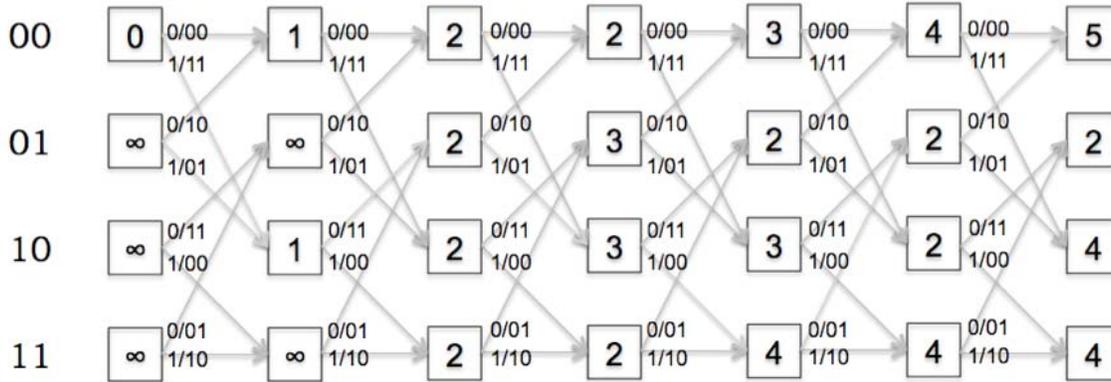
D. If the transmitted message starts with the bits "01011", what is the sequence of bits produced by the convolutional encoder?

Show Answer

The receiver determines the most-likely transmitted message by using the Viterbi algorithm to process the (possibly corrupted) received parity bits. The path metric trellis generated from a particular set of received parity bits is shown below. The boxes in the trellis contain the minimum path metric as computed by the

Viterbi algorithm.

Time step	1	2	3	4	5	6
Received	01	01	00	01	01	11



E. Referring to the trellis above, what is the receiver's estimate of the most-likely transmitter state after processing the bits received at time step 6?

Show Answer

F. Referring to the trellis above, show the most-likely path through the trellis by placing a circle around the appropriate state box at each time step and darkening the appropriate arcs. What is the receiver's estimate of the most-likely transmitted message?

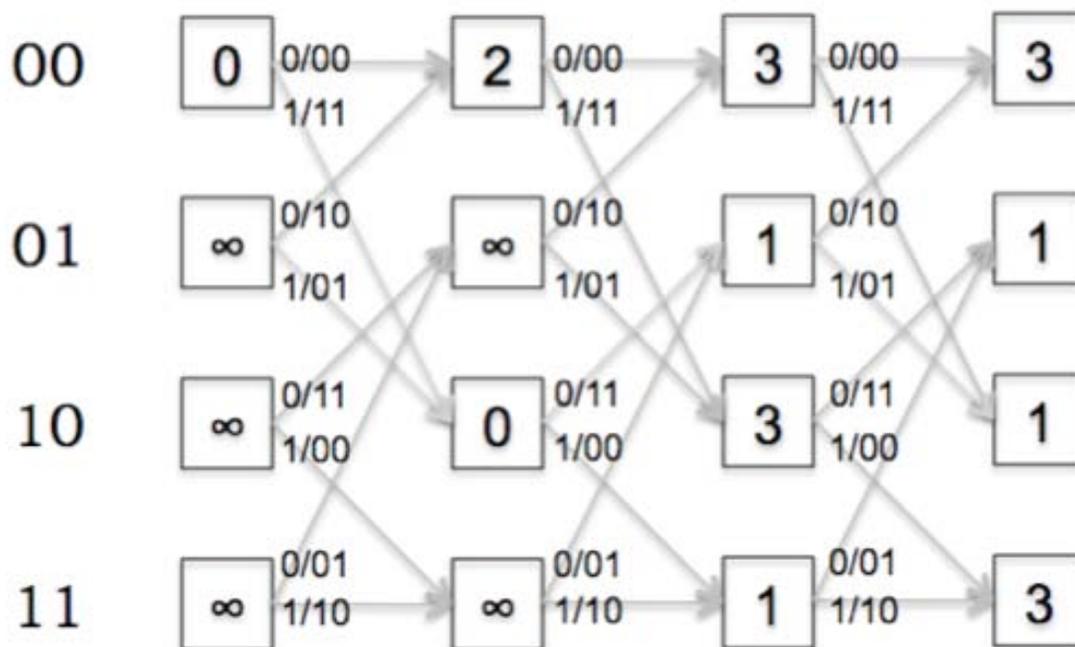
Show Answer

G. Referring to the trellis above, and given the receiver's estimate of the most-likely transmitted message, at what time step(s) were errors detected by the receiver? Briefly explain your reasoning.

Show Answer

H. Now consider the path metric trellis generated from a *different* set of received parity bits.

Time step	1	2	3
Received	??	??	??



Referring to the trellis above, determine which pair(s) of parity bits could have been received at time steps 1, 2 and 3. Briefly explain your reasoning.

[Show Answer](#)

Problem 11. Consider a binary convolutional code specified by the generators (1011, 1101, 1111).

- A. What are the values of
 - a. constraint length of the code
 - b. rate of the code
 - c. number of states at each time step of the trellis
 - d. number of branches transitioning into each state
 - e. number of branches transitioning out of each state
 - f. number of expected parity bits on each branch

[Show Answer](#)

A 10000-bit message is encoded with the above code and transmitted over a noisy channel. During Viterbi decoding at the receiver, the state 010 had the lowest path metric (a value of 621) in the final time step, and the survivor path from that state was traced back to recover the original message.

- B. What is the likely number of bit errors that are corrected by the decoder? How many errors are likely left uncorrected in the decoded message?

[Show Answer](#)

- C. If you are told that the decoded message had no uncorrected errors, can you guess the approximate number of bit errors that would have occurred had the 10000 bit message been transmitted without any coding on the same channel?

Show Answer

- D. From knowing the final state of the trellis (010, as given above), can you infer what the last bit of the original message was? What about the last-but-one bit? The last 4 bits?

Show Answer

Consider a transition branch between two states on the trellis that has 000 as the expected set of parity bits. Assume that 0V and 1V are used as the signaling voltages to transmit a 0 and 1 respectively, and 0.5V is used as the digitization threshold.

- E. Assuming hard decision decoding, which of the two set of received voltages will be considered more likely to correspond to the expected parity bits on the transition: (0V, 0.501V, 0.501V) or (0V, 0V, 0.9V)? What if one is using soft decision decoding?

Show Answer

Problem 12. Indicate whether each of the statements below is true or false, and a brief reason why you think so.

- A. If the number states in the trellis of a convolutional code is S , then the number of survivor paths at any point of time is S . Remember that if there is "tie" between two incoming branches (i.e., they both result in the same path metric), we arbitrarily choose only one as the predecessor.

Show Answer

The path metric of a state s_1 in the trellis indicates the number of residual uncorrected errors left along the trellis path from the start state to s_1 .

Show Answer

- B. Among the survivor paths left at any point during the decoding, no two can be leaving the same state at any stage of the trellis.

Show Answer

- C. Among the survivor paths left at any point during the decoding, no two can be entering the same state at any stage of the trellis. Remember that if there is "tie" between two incoming branches (i.e., they both result in the same path metric), we arbitrarily choose only one as the predecessor.

Show Answer

- D. For a given state machine of a convolutional code, a particular input message bit stream always produces the same output parity bits.

Show Answer

Problem 13. Consider a convolution code with two generator polynomials: $G_0=101$ and $G_1=110$.

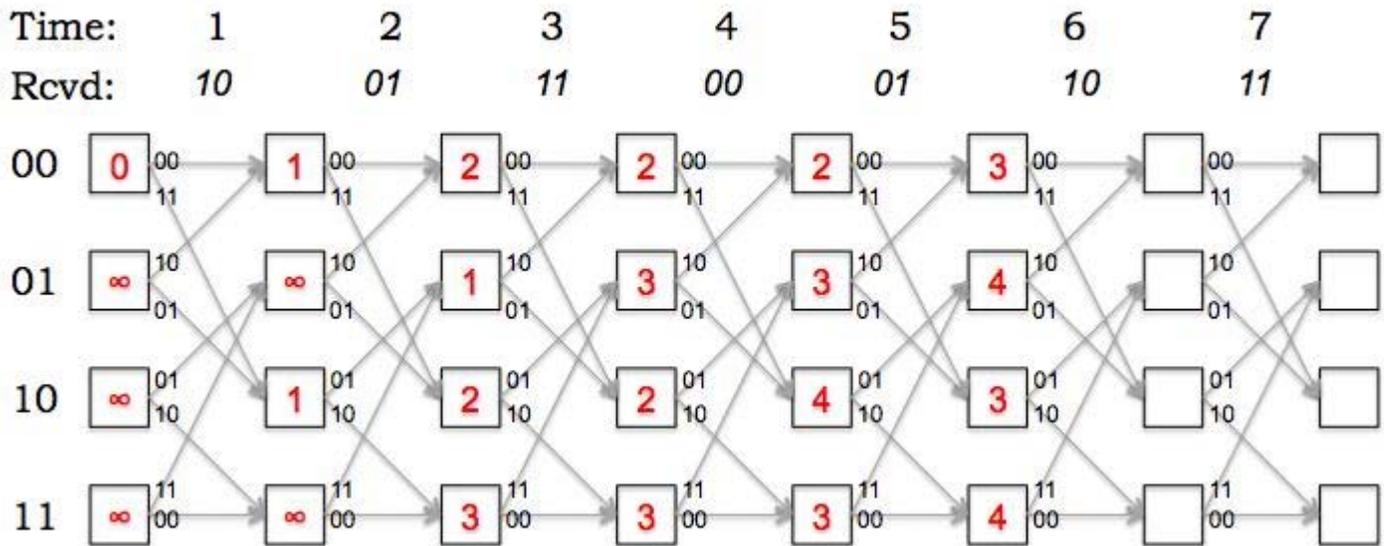
A. What is code rate r and constraint length k for this code?

Show Answer

B. Draw the state transition diagram for a transmitter that uses this convolutional code. The states should be labeled with the binary string $x_{n-1} \dots x_{n-k+1}$ and the arcs labeled with x_n/p_0p_1 where $x[n]$ is the next message bit and p_0 and p_1 are the two parity bits computed from G_0 and G_1 respectively.

Show Answer

The figure below is a snapshot of the decoding trellis showing a particular state of a maximum likelihood decoder implemented using the Viterbi algorithm. The labels in the boxes show the path metrics computed for each state after receiving the incoming parity bits at time t . The labels on the arcs show the expected parity bits for each transition; the actual received bits at each time are shown above the trellis.



C. Fill in the path metrics in the empty boxes in the diagram above (corresponding to the Viterbi calculations for times 6 and 7).

Show Answer

D. Based on the updated trellis, what is the most-likely final state of the transmitter? How many errors were detected along the most-likely path to the most-likely final state?

Show Answer

E. What's the most-likely path through the trellis (i.e., what's the most-likely sequence of states for the transmitter)? What's the decoded message?

Show Answer

F. Based on your choice of the most-likely path through the trellis, at what times did the errors occur?

Show Answer

MIT OpenCourseWare
<http://ocw.mit.edu>

6.02 Introduction to EECS II: Digital Communication Systems
Fall 2012

For information about citing these materials or our Terms of Use, visit: <http://ocw.mit.edu/terms>.