

INTRODUCTION TO EECS II
DIGITAL COMMUNICATION SYSTEMS

**6.02 Fall 2012
 Lecture #23**

A Brief History of the Internet

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The Dawn of Packet Switching

ARPA: 1957, in response to Sputnik
 Paul Baran (RAND Corp)

- Early 1960s: New approaches for survivable comms systems; “hot potato routing” and decentralized architecture, paper on packet switching over *digital* links

Donald Davies (UK), early 1960s

- Coins the term “packet”

Len Kleinrock (MIT thesis): “Information flow in large communication nets”, 1961

J. Licklider & W. Clark (MIT), *On-line Man Computer Communication* (1962) & Licklider’s vision of a “galactic network”

L. Roberts (MIT then ARPA), first ARPANET plan for time-sharing remote computers

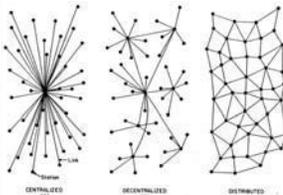


FIG. 1 - Centralized, Decentralized and Distributed Networks

RAND Corporation, *On Distributed Communications: Introduction to Distributed Communications Networks*, RM-3420-PR, 1964. Reprinted with permission.

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ARPANET



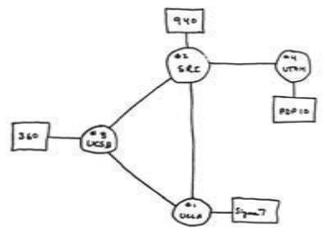
BBN team that implemented the *interface message processor (IMP)*

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- 1967: Connect computers at key research sites across the US using telephone lines
- Interface Message Processors (IMP) ARPA contract to BBN
- Sen. Ted Kennedy sends a somewhat confused telegram to BBN on winning the contract
 “Congratulations ... on **interfaith** message processor”

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Initial Baby Steps



THE ARPA NETWORK
 DEC 1969
 4 NODES

FIGURE 6.2 Drawing of 4 Node Network (Courtesy of Alex McKenzie)

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In the Beginning...

- Kleinrock’s group at UCLA tried to log on to SRI computer: His recollection of the event...
- “We set up a telephone connection between us and the guys at SRI...
- We typed the L and we asked on the phone...
 - “Do you see the L?”
 - “Yes, we see the L,” came the response
- We typed the O, and we asked...
 - “Do you see the O?”
 - “Yes, we see the O.”
- Then we typed the G...
 - ...and the system crashed!



Image in the public domain. Source: [Wikipedia](https://en.wikipedia.org/wiki/ARPANET).

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September 1971

1970, ARPANET hosts start using NCP; first two cross-country lines (BBN-UCLA and MIT-Utah)
 “Hostile overlay” atop telephone network
 Ran a distance-vector routing protocol



MAP 4 September 1971

Image in the public domain, from the *ARPANET Completion Report*, January 1978.

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1970s: Packet networks → Internetworking

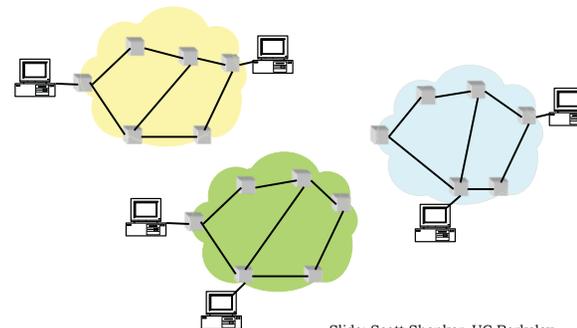
- 1972: successful ARPANET demo at conference (except it failed when demo'd to skeptics from AT&T!)
- 1972: modified ARPANET email program
- 1972: CYCLADES network (Louis Pouzin et al.): best-effort “datagrams”; *sliding window* protocol; distance-vector routing; time sync
- 1973: Ethernet (MAC protocol inspired by Aloha -- CSMA)
- 1973-74: Xerox PUP (used distance-vector protocol)
- 1973: ARPANET becomes international
- 1973-75: Internetworking effort (Cerf, Kahn, et al.)
 - Developed TCP and IP (originally intertwined) – TCP uses *sliding window*

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The Problem

- Many different packet-switching networks
- Only nodes on the same network could communicate



Slide: Scott Shenker, UC Berkeley

Courtesy of Scott Shenker. Used with permission.

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Kahn's Rules for Interconnection

- Each network is independent and must not be required to change
- Best-effort communication
- Boxes (then called *gateways*) connect networks
- No global control at operations level (why?)

Cerf RFC 968

Twas the night before start-up and all through the net,
not a packet was moving; no bit nor octet.
The engineers rattled their cards in despair,
hoping a bad chip would blow with a flare.
The salesmen were nestled all snug in their beds,
while visions of data nets danced in their heads.
And I with my datascopce tracings and dumps
prepared for some pretty bad bruises and lumps.
When out in the hall there arose such a clatter,
I sprang from my desk to see what was the matter.

There stood at the threshold with PC in tow,
An ARPANET hacker, all ready to go.
I could see from the creases that covered his brow,
he'd conquer the crisis confronting him now.
More rapid than eagles, he checked each alarm
and scrutinized each for its potential harm.

Courtesy of Vint Cerf. Used with permission.

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Solution

Courtesy of Scott Shenker. Used with permission.

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Slide: Scott Shenker, UC Berkeley
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The Internetworking Vision

- Bob Kahn & Vint Cerf imagined there would be only a few networks and thus only a few gateways
 - “*The choice for network identification (8 bits) allows up to 256 distinct networks. This size seems sufficient for the foreseeable future.*”
 - They were a little wrong!
- Gateways would “translate” between networks
 - Evolved in the 1974 Cerf/Kahn paper as a universal network layer, later called the *Internet Protocol*, or IP
 - We now think of it as all routers supporting IP

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Handling Heterogeneity

- Make it very easy to be a node or link on the network (best-effort)
- Universal *network layer*: standardize addressing and forwarding
- Switches maintain *no per-connection state* on behalf of end points
- Original addressing model (then called a “TCP address” because IP and TCP hadn’t been split into different layers):

8	16
NETWORK	TCP IDENTIFIER
TCP address	

Image by MIT OpenCourseWare.

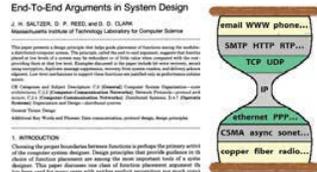
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1970s: Internetworking

- 1978: **Layering!** TCP and IP split; TCP at end points, IP in the network
- IP network layer: simple best-effort delivery
- In retrospect: Packet switching (& TCP/IP) won because it is good enough for almost every application (though optimal for nothing!)
- Competitor to TCP/IP: ISO, standardizing 7-layer *OSI stack*

Classic Internet layering "hourglass" model



End-To-End Arguments in System Design

J. H. KALCER, D. P. FRED, and D. D. CLARK
Massachusetts Institute of Technology, Laboratory for Computer Science

The network is a design space with many degrees of freedom. It is a complex, multi-layered system. The network is a design space with many degrees of freedom. It is a complex, multi-layered system. The network is a design space with many degrees of freedom. It is a complex, multi-layered system.

1. INTRODUCTION
Choosing the proper layering between functions is perhaps the primary activity of the system designer. Design principles that provide guidance in this activity are presented in this paper. The design principles are based on the end-to-end argument. This paper discusses the role of function placement arguments in the design of a network. The design principles are based on the end-to-end argument. This paper discusses the role of function placement arguments in the design of a network.

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Most Useful Lessons

*One should architect systems for flexibility – you'll almost never know what apps make it succeed.
(Even if it means sacrificing some performance!)*

Il semble que la perfection soit atteinte non quand il n'y a plus rien à ajouter, mais quand il n'y a plus rien à retrancher.
Perfection is achieved, not when there is nothing more to add, but when there is nothing left to take away

-- Antoine de Saint-Exupery

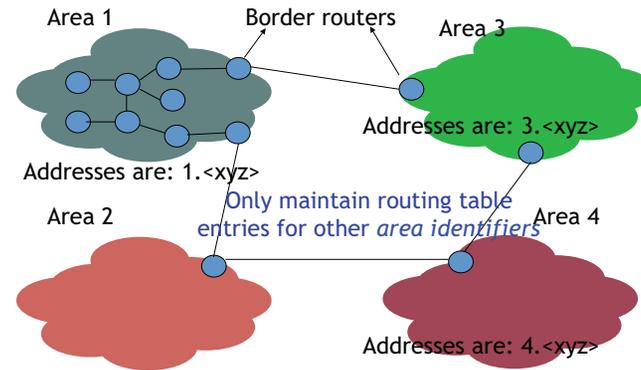
Or,

When in doubt, leave it out

1980s: Handling Growth with Topological Addressing

- 1978-79: ARPANET moves to *link-state routing*
- Per-node routing entries don't scale well
- Solution: Organize network hierarchically
 - Into "areas" or "domains"
 - Similar to how the postal system works
 - Hide detailed information about remote areas
- For this approach to work, node addresses must be *topological*
 - Address should tell network *where* in the network the node is
 - I.e., address is a *location* in the network
- Three classes of addresses in the 80s: "Class A", "Class B", and "Class C"
 - Not used any more, though the dotted decimal notation of IPv4 addresses makes it look like the dots matter

Ideal Case: Classic "Area Routing"



1980s: Rapid Growth

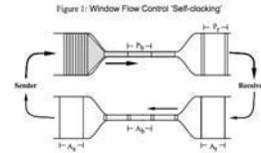
- 1981-89: Dave Clark of MIT is Internet's "Chief Architect"
 - Co-author of the end-to-end arguments (w/ Saltzer/Reed)
 - Ensures consistency of design and vision
 - "We reject kings, presidents, and voting. We believe in rough consensus and running code."
- 1982: US DoD standardizes on TCP/IP
 - Berkeley's computer systems research group produces BSD & sockets
- 1983: MIT Project Athena - large-scale campus-area networking
- 1984: Domain Name System (DNS) introduced
- 1985: NSFNet picks TCP/IP as standard

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Growth Problems: Congestion

- 1986: Congestion collapse episodes on the Internet
 - Problems with bad timeout settings
 - Window size not appropriate for network state
 - Athena network file system congestion problems (bad timeout settings)
- Congestion avoidance and control
 - RTT estimation using EWMA + new timeout method
 - TCP congestion control by Van Jacobson (concurrent work on DECBit scheme by Ramakrishnan & Jain)
 - Adapt the window size to congestion: If congested, decrease window; else increase. Use exponential back-offs on timeouts
 - By the end of the 1980s, essentially all running TCPs had congestion control



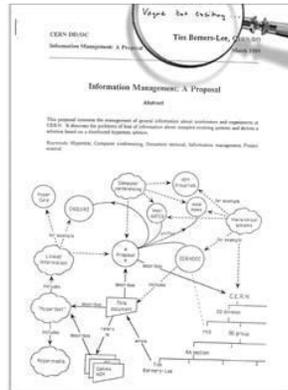
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1990s

- 1990: no more ARPANET
- 1991: Tim Berners-Lee releases "WorldWideWeb"
- Mid-1990s: NSFNet backbone ends
 - Commercial ISPs take off
- "Classless" addressing for scale
 - And the rise of NATs
- BGP4: Path vector protocol between competing ISPs, who must yet cooperate



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1990s: Handling Growth with CIDR IPv4 Addresses & Address Prefixes

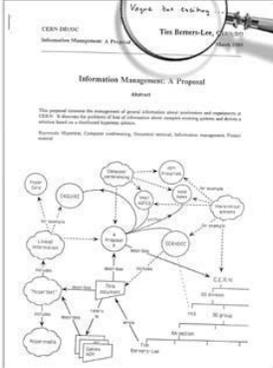
- 18.31.0.82 is actually the 32 bit string 0001001000111110000000001010010
- Routers have forwarding table entries corresponding to an address **prefix** (a range of addrs w/ common prefix bitstring)
- 18.0.0.0/8 stands for all IP addresses in the range 00010010 00...0 to 00010010 11...1 (i.e., 2^{24} addresses of the form 00010010*)
- 18.31.0.0/17 stands for a range of 2^{15} consecutive IP addresses of the form 00010010001111100* (1st 17 bits are the same for each address in that range)
- Hence, *subnetworks* may be of size 1, 2, 4, 8, ... (maxing out at 2^{24} usually), and may be recursively divided further
- Forwarding uses **longest prefix match**
 - At each router, routes are of the form "For this range of addresses, use this route"
 - Pick the route that has the longest matching prefix w/ dest addr

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- BGP4: *Path vector protocol* between competing ISPs, who must yet cooperate
- 1991-1994: IPng & IPv6 design starts
- 1993: search engines (Excite)
- Mid-1990s: E-commerce starts
- 1998: Google reinvents search
- 1998: Content distribution networks like Akamai
- 1996-2001: .com bubble starts & bursts



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2000s: The Internet Matures (Top 5 List)

- 2000-2001: .com bust
 - And 9/11 happens
 - Power of content distribution networks to handle load
- 1. The rise of peer-to-peer networks
 - Gnutella, Freenet, distributed hash tables (e.g., Chord), BitTorrent, and of course, Napster
- 2. Security threats and defenses
 - 2000: Large-scale distributed denial-of-service (DDoS) attacks start
 - 2003: SQL slammer worm
 - Spam → phishing and pharming → complex ecosystem
 - Route hijacking by errors or malice
- 3. User-generated content & social networks
 - Blogs, Youtube, Facebook, and Twitter (UGC-meets-social)
- 4. The rise of wireless and mobile data
- 5. Cloud computing and large-scale datacenters (Amazon, Google, Microsoft, Facebook, etc.)

- *Almost everything* moves to the Internet: telephony, video, entertainment

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2010s: The Decade Ahead

- Even more massive growth (largely from video, entertainment, and collaboration) & internationalization
- Combating complexity: new methods to make things simpler (“software-defined networks”)
- Wireless: the spectrum crisis, coping with mobility & variability
- Physical embedding & embodiment: sensors & actuators over the network, mobile robots and autonomous agents, vehicles, embedded devices
- Network security & privacy; censorship

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6.02 in One Slide



How to design digital communication networks.
 Three layers of abstraction: bits, signals, packets.
 A unique storyline: vertical study across all layers

Bits: Binary representation. Compression (source coding). Bit errors and error correction codes (channel coding)

Signals: Noise. LTI models. Frequency-domain analysis. Modulation & demodulation.

Packets: MAC protocols for shared media. Packet-switching & queues. Routing protocols. Reliable transport.

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What Next?

- Many UROP opportunities!
- Networks and computer systems
 - 6.033 (computer systems), 6.829 (computer networks), 6.824 (distributed systems), 6.263 (analysis of networks), 6.266 (network algorithms)
- Security
 - 6.857 (computer and network security), 6.858 (computer systems security)
- Signal processing & digital communications
 - 6.003 (signals and systems), 6.011 (communications, control, and signal processing)
- Advanced communication & information theory
 - 6.450 & 6.451 (digital communications), 6.441 (info theory)

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Thank you!

- Lectures
 - George Verghese
 - Hari Balakrishnan
- Recitations
 - Yury Polyanskiy
 - Jacob White
 - Victor Zue
- TAs
 - Rui Hu
 - Shao-Lun Huang
 - Ruben Madrigal
 - Kyu Seob Kim
 - Eduardo Sverdlin-Lisker
 - Cassandra Xia

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