

1.204 Lecture 20

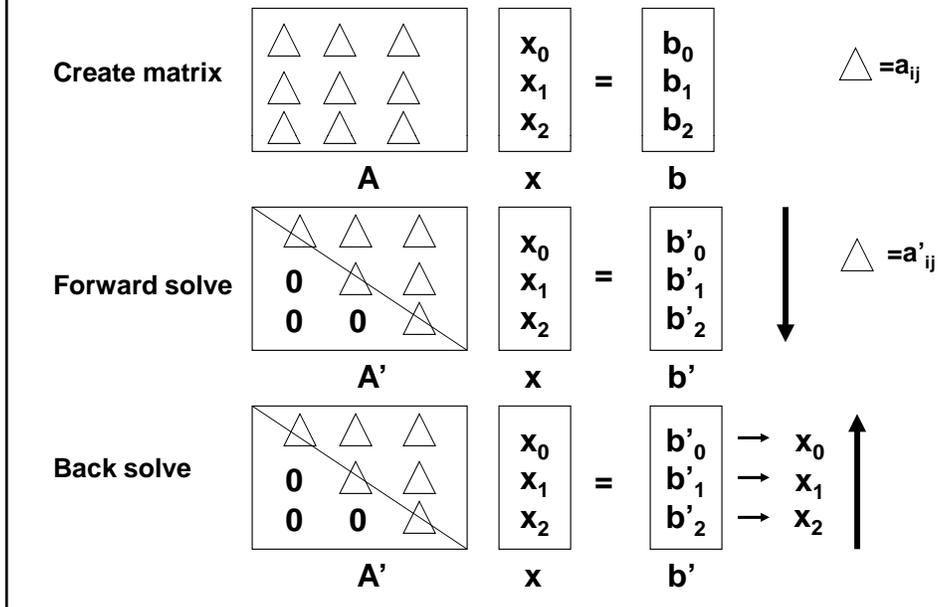
Linear systems:
Gaussian elimination
LU decomposition

Systems of Linear Equations

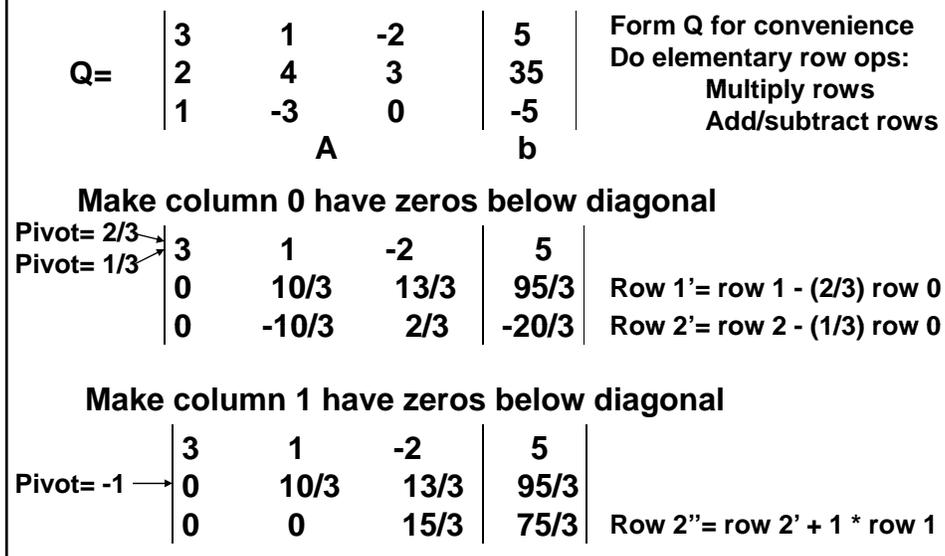
$$\begin{aligned} 3x_0 + x_1 - 2x_2 &= 5 \\ 2x_0 + 4x_1 + 3x_2 &= 35 \\ x_0 - 3x_1 &= -5 \end{aligned}$$

$$\begin{array}{ccc|ccc} 3 & 1 & -2 & x_0 & & 5 \\ 2 & 4 & 3 & x_1 & & 35 \\ 1 & -3 & 0 & x_2 & & -5 \\ \hline & A & & x & = & b \\ & 3 \times 3 & & 3 \times 1 & & 3 \times 1 \end{array}$$

Algorithm to Solve Linear System



Gaussian Elimination: Forward Solve



Gaussian Elimination: Back Solve

3	1	-2	5
0	10/3	13/3	95/3
0	0	15/3	75/3

$$(15/3)x_2 = (75/3)$$

$$x_2 = 5$$

3	1	-2	5
0	10/3	13/3	95/3
0	0	15/3	75/3

$$(10/3)x_1 + (13/3)*5 = (95/3) \quad x_1 = 3$$

3	1	-2	5
0	10/3	13/3	95/3
0	0	15/3	75/3

$$3x_0 + 1*3 - 2*5 = 5$$

$$x_0 = 4$$

A Complication

0	1	-2	5
2	4	3	35
1	-3	0	-5

$$\text{Row 1}' = \text{row 1} - (2/0) \text{row 0}$$

Exchange rows: put largest pivot element in row:

2	4	3	35
0	1	-2	5
1	-3	0	-5

Do this as we process each column.

If there is no nonzero element in a column, matrix is not full rank.

Gaussian Elimination

```
public class Gauss {
    public static double[] gaussian(double[][] a, double[] b) {
        int n = a.length;           // Number of unknowns
        double[][] q = new double[n][n + 1];

        for (int i = 0; i < n; i++) {
            for (int j = 0; j < n; j++) // Form q matrix
                q[i][j] = a[i][j];
            q[i][n] = b[i];
        }
        forward_solve(q);           // Do Gaussian elimination
        back_solve(q);              // Perform back substitution

        double[] x = new double[n]; // Extract column n of q,
        for (int i = 0; i < n; i++) // which contains the solution x
            x[i] = q[i][n];
        return x;
    }
}
```

Forward Solve

```
private static void forward_solve(double[][] q) {
    int n = q.length;
    int m = q[0].length;

    for (int i = 0; i < n; i++) { // Find row w/max element in this
        int maxRow = i;         // column, at or below diagonal
        for (int k = i + 1; k < n; k++)
            if (Math.abs(q[k][i]) > Math.abs(q[maxRow][i]))
                maxRow = k;

        if (maxRow != i) // If row not current row, swap
            for (int j = i; j < m; j++) {
                double t = q[i][j];
                q[i][j] = q[maxRow][j];
                q[maxRow][j] = t;
            }

        for (int j = i + 1; j < n; j++) { // Calculate pivot ratio
            double pivot = q[j][i] / q[i][i];
            for (int k = i; k < m; k++) // Pivot operation itself
                q[j][k] -= q[i][k] * pivot;
        }
    }
}
```

Back Substitution

```
private static void back_solve(double[][] q) {
    int n = q.length;
    int m = q[0].length;
    for (int p = n; p < m; p++) {           // Loop over p columns
        for (int j = n - 1; j >= 0; j--) { // Start at last row
            double t = 0.0;                // t- temporary
            for (int k = j + 1; k < n; k++)
                t += q[j][k] * q[k][p];
            q[j][p] = (q[j][p] - t) / q[j][j];
        }
    }
}
```

Variations

Multiple right hand sides: augment Q, solve all eqns at once

$$\left| \begin{array}{ccc|c|c|c} 3 & 1 & -2 & 5 & 7 & 87 \\ 2 & 4 & 3 & 35 & 75 & -1 \\ 1 & -3 & 0 & -5 & 38 & 52 \end{array} \right|$$

Matrix inversion (rarely done in practice)

$$\left[\begin{array}{ccc|ccc} 3 & 1 & -2 & 1 & 0 & 0 \\ 2 & 4 & 3 & 0 & 1 & 0 \\ 1 & -3 & 0 & 0 & 0 & 1 \end{array} \right] \longrightarrow \left[\begin{array}{ccc|ccc} \# & \# & \# & @ & @ & @ \\ 0 & \# & \# & @ & @ & @ \\ 0 & 0 & \# & @ & @ & @ \end{array} \right]$$

$\underbrace{\hspace{10em}}_Q$
 $A \bullet ? = I$
 A^{-1}

$? = A^{-1}$

Invert

```

public static double[][] invert(double[][] a) {
    int n = a.length;           // Number of unknowns
    double[][] q = new double[n][n+n];

    for (int i = 0; i < n; i++)
        for (int j = 0; j < n; j++) // Form q matrix
            q[i][j] = a[i][j];

    // Form identity matrix in right half of q
    for (int i = 0; i < n; i++)
        q[i][n+i] = 1.0;

    forward_solve(q);           // Do Gaussian elimination
    back_solve(q);              // Perform back substitution

    double[][] x = new double[n][n]; // Extract R half of q
    for (int i = 0; i < n; i++) // which contains inverse
        for (int j = 0; j < n; j++)
            x[i][j] = q[i][j+n];
    return x;
}
// Method multiply() in download
// Example use in GaussTest in download

```

LU decomposition

- We can write matrix **A** as the product of two matrices **L** and **U**:

$$\begin{bmatrix} l_{00} & 0 & 0 & 0 \\ l_{10} & l_{11} & 0 & 0 \\ l_{20} & l_{21} & l_{22} & 0 \\ l_{30} & l_{31} & l_{32} & l_{33} \end{bmatrix} \cdot \begin{bmatrix} u_{00} & u_{01} & u_{02} & u_{03} \\ 0 & u_{11} & u_{12} & u_{13} \\ 0 & 0 & u_{22} & u_{23} \\ 0 & 0 & 0 & u_{33} \end{bmatrix} = \begin{bmatrix} a_{00} & a_{01} & a_{02} & a_{03} \\ a_{10} & a_{11} & a_{12} & a_{13} \\ a_{20} & a_{21} & a_{22} & a_{23} \\ a_{30} & a_{31} & a_{32} & a_{33} \end{bmatrix}$$

- We can solve

$$A \cdot x = (L \cdot U) \cdot x = L \cdot (U \cdot x) = b$$

by first solving for a vector **y**

$$L \cdot y = b$$

and then solving

$$U \cdot x = y$$

Why? Solving each is trivial: forward, back substitution

Why and How

- This is perhaps twice as fast as Gaussian elimination (count steps)
- L and U do not depend on b, so we can solve as many right hand sides as we wish
- **How: Crout's method**
 - We can decompose matrix A into matrices L and U by arranging the equations in a given order
 - The rearrangement is subtle; we don't cover it in class since you'll never need to implement or modify it
- Java implementation is on the Web site, based on Press et al, *Numerical Recipes*

Class LU

- **Constructor: LU(double[][] a)**
 - Stores LU decomposition in a single matrix
 - All $l_{ij} = 1.0$ in matrix L
 - We store all u elements and all non-diagonal l elements in LU
- **Methods:**
 - public double[] solve(double[] b)
 - public double[][] solve(double[][] b)
 - public double[][] inverse()
 - public double determinant()
 - public double[] improve(double[] b, double[] x)
- **See download for code and LUTest class for examples of usage**
 - You can use it as a 'black box'
 - Use this in preference to class Gauss

Other linear system algorithms

- **Banded matrices**
- **Sparse matrices**
- **Singular value decompositions (SVD)**
 - Should be used in least squares computations
- **Cholesky decomposition ($A = L L^T$)**
 - Square, symmetric, positive definite matrices
 - Used in econometrics
- **And others...**
 - Almost all are based on pivot operations

Linear system model: Rail performance

- **Compute running time, including delays, for trains on a single track railroad**
 - Traffic in both directions: east and west
 - Three types of train (6 classes of train, including direction)

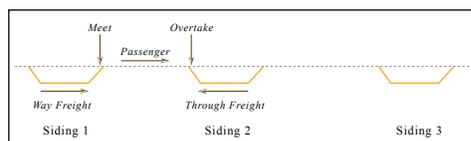


Figure by MIT OpenCourseWare.

Class	Description	Velocity
0	WB way freight	-25
1	WB thru freight	-50
2	WB passenger	-80
3	EB passenger	80
4	EB thru freight	50
5	EB way freight	25

From E.R. Petersen

Rail performance, p.2

- **Priorities used to model meets and overtakes**
 - Meets occur when trains travel in opposite directions and one must take siding and wait until other passes
 - Overtakes occur when trains traveling in same direction interact, and slower one takes siding to let faster one go by
 - Assume all sidings are long enough
- **We want to model the running time for each class of train over a segment of railroad, as a function of:**
 - Number of trains of each class (type, direction)
 - Speed of trains
 - Number of sidings
 - Priorities, and other, less important variables
- **Our linear model will give nonlinear performance behavior!**

Delay matrix D

- **Matrix D gives average delay for each interaction (meet or overtake) between two classes of train**
 - We will then multiply this by the expected number of interactions, to get total delay
- **Delay matrix D has coefficients D_{ij} :**

$$D_{ij} = p_{ij}S_i + \frac{60p_{ij}^2}{2(b+1)} \left| \frac{d}{v_i} - \frac{d}{v_j} \right|$$

- D_{ij} = Expected delay to train i due to train j, in minutes
- S_i = Time to take siding for train i, in minutes (5 minutes)
- p_{ij} = Relative number of times train i waits for train j ($0 \leq p_{ij} \leq 1$)
- b = Number of sidings (19)
- D = Distance of railroad segment being modeled, in miles (400 mi)
- v_i = Free running velocity for train i, in miles per hour

Delay matrix D

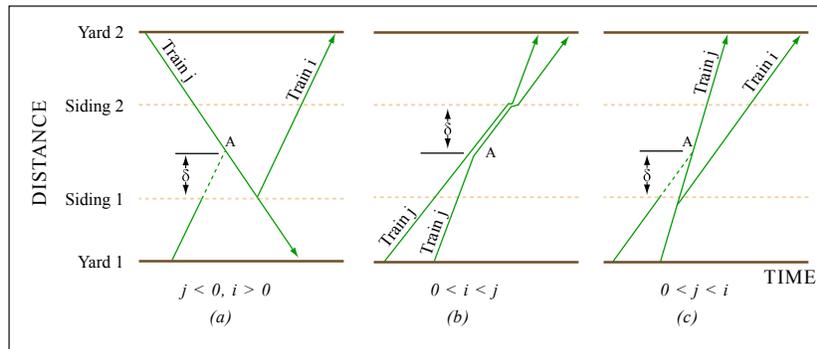


Figure by MIT OpenCourseWare.

Probability, delay matrices

Prob(delay)	WB way	WB thru	WB pass	EB pass	EB thru	EB way
WB way	0	0.7	0.9	1	0.7	0.5
WB thru	0.3	0	0.7	0.7	0.5	0.3
WB pass	0.1	0.3	0	0.5	0.3	0
EB pass	0	0.3	0.5	0	0.3	0.1
EB thru	0.3	0.5	0.7	0.7	0	0.3
EB way	0.5	0.7	1	0.9	0.7	0

Delay	WB way	WB thru	WB pass	EB pass	EB thru	EB way
WB way	0	9.4	17.9	36.5	21.1	14.5
WB thru	2.6	0	5.7	13.1	8.5	4.7
WB pass	0.7	1.9	0	6.3	3.3	0
EB pass	0	3.3	6.3	0	1.9	0.7
EB thru	4.7	8.5	13.1	5.7	0	2.6
EB way	14.5	21.1	36.5	17.9	9.4	0

Derivation of linear system

- Let
 - W_i = average time for train of class i, including delays
 - T_i = free running time for train of class i (input)
 - D_{ij} = delays due to meets and overtakes to train of class i due to trains of class j
 - M_{ij} = number of meets and overtakes between train of class i and trains of class j

- Average time for train to travel across segment:

$$W_i = T_i + \sum_j (D_{ij} \cdot M_{ij})$$

- Interactions between train i and trains of class j:

$$M_{ij} = N_j (W_j + W_i) / 1440$$

- M_{ij} = number of trains/day in class j times fraction of day that train of class i can interact with trains of class j
 - If EB train takes 12 hours (720 minutes) to cover line, as does WB, it will meet every WB train that operates that day
 - If EB train took only 6 hours, it would have half the interactions

Derivation of linear system, p.2

$$W_i = T_i + \sum_j (D_{ij} \cdot M_{ij})$$

$$M_{ij} = N_j (W_j + W_i) / 1440$$

- The rest is algebra, to write the two equations above in the form $AW = T$, with W the unknown (“x”)

$$W_i = T_i + \sum_j \frac{D_{ij} \cdot N_j}{1440} (W_j + W_i)$$

Collect W_i terms on left side of equation :

$$W_i - \sum_j \frac{D_{ij} \cdot N_j}{1440} (W_j + W_i) = T_i$$

Rearrange terms to separate W_i and W_j terms :

$$\left(1 - \sum_j \frac{D_{ij} \cdot N_j}{1440}\right) \cdot W_i - \sum_j \frac{D_{ij} \cdot N_j}{1440} \cdot W_j = T_i$$

The W_i coefficient is the diagonal; the W_j coefficients are the off - diagonal elements in matrix A

$$\begin{array}{ccc|c|c} a_{00} & a_{01} & a_{02} & w_0 & t_0 \\ a_{10} & a_{11} & a_{12} & w_1 & t_1 \\ a_{20} & a_{21} & a_{22} & w_2 & t_2 \end{array}$$

- There is a +/- sign convention handled by a matrix C (multiplies D_{ij}):
 - $c[i][j] = -1$ if $j < i \leq 2$ or $(3 \leq i < j)$, 0 otherwise (with 6 train classes)

Data members, constructor

```
public class RailDelay {
    private int n; // Number of train classes
    private int d; // Distance in miles
    private int b; // Number of sidings
    private double[][] p; // Probability of delay in interaction
    private double[] s; // Time to take siding, by train class
    private double[] v; // Velocity by train class
    private double[][] c; // Matrix that indicates if interaction
    // is pass or meet, by train classes involved
    private int[] nTrain; // Number of trains by class, per day

    public RailDelay(String filename) {
        // Constructor reads inputs from file, sets data members
    }
}
```

getDelay()

```
public double[] getDelay() {
    double[][] dm= new double[n][n]; // Delay matrix D
    for (int i= 0; i < n; i++)
        for (int j= 0; j < n; j++) {
            dm[i][j]= p[i][j]*s[i] + 60.0*p[i][j]*p[i][j] *
                Math.abs(d/v[i]- d/v[j])/(2.0*(b+1)); }
    double[] t= new double[n]; // Free running time T
    for (int i= 0; i < n; i++)
        t[i]= d*60.0/v[i];
    double[][] a= new double[n][n]; // Total delay matrix A
    for (int i= 0; i < n; i++) {
        double delay= 0.0;
        for (int j= 0; j < n; j++) {
            if (i != j) {
                a[i][j]= dm[i][j]*nTrain[j]*c[i][j]/1440;
                delay += a[i][j];
            }
        }
        a[i][i]= 1.0 - delay; }
    double[] w= Gauss.gaussian(a, t);
    return w;
}
```

main(), sample output

```
public static void main(String[] args) {
    RailDelay r= new RailDelay("src/linear/rail.txt");
    double[] w= r.getDelay();    // Gets output w
    int n= w.length;
    System.out.println("i Act time Free time");
    for (int i= 0; i < n; i++)
        System.out.printf("%d %8.1f %8.1f \n", i,
            Math.abs(w[i]), Math.abs(r.getFreeTime(i)));
    System.out.println();
}
// Sample output: 3 wayfreight, 4 thru freight, 2 passenger
i Act time Free time
0 1265.2 960.0
1 544.9 480.0
2 315.8 300.0
3 315.8 300.0
4 544.9 480.0
5 1265.2 960.0
```

Rail performance estimate

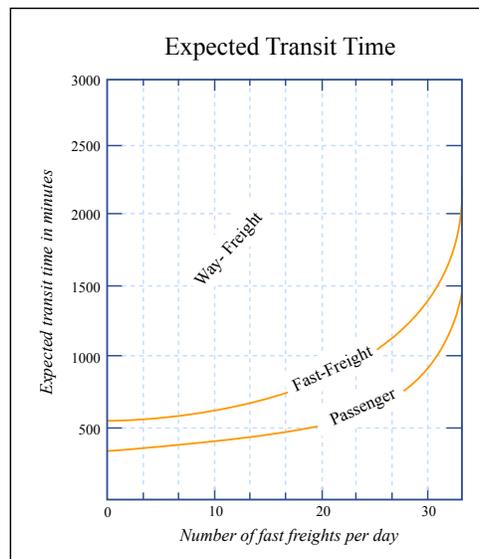


Figure by MIT OpenCourseWare.

Summary

- **Linear models are a reasonable starting point in many cases to understand complex systems**
 - Writing down equations to model a system analytically or through solving linear or nonlinear systems is often a viable option
 - Linear models can produce nonlinear behavior
 - In the rail example, this is more intuitive (for some of us) and more robust than simulation
- **I used essentially the rail analysis in a Vermont Act 250 expert witness case**
 - Traffic impacts on a neighborhood from a large development
 - Narrow road with parking on both sides
 - Number of “meets” between cars would increase very sharply
 - Project application was denied
- **We’ll use linear systems as a “subproblem” in next lecture**

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